

Számítógépes számelmélet

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Ezek a programok csak szemléltetésre szolgálnak

- ▶ 1. A prímek eloszlása, szitálás
- ▶ 2. Egyszerű faktorizálási módszerek
- ▶ 3. Egyszerű prímtesztelési módszerek
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- ▶ 10. Faktorizálás elliptikus görbékkel
- ▼ 11. Prímteszt elliptikus görbékkel

```
> restart; with(numtheory);  
[Glgcd, bigomega, cfrac, cfracpol, cyclotomic, divisors, factorEQ, factorset, (11.1)  
fermat, imagunit, index, integral_basis, invcfrac, invphi, issqrfree, jacobi,  
kronecker,  $\lambda$ , legendre, mcombine, mersenne, migcdex, minkowski,  
mipolys, mlog, mobius, mroot, msqrt, nearestp, nthconver, nthdenom,  
nthnumer, nthpow, order, pdexpand,  $\phi$ ,  $\pi$ , pprimroot, primroot, quadres,  
rootsunity, safeprime,  $\sigma$ , sq2factor, sum2sqr,  $\tau$ , thue]
```

▼ 11.1. Tétel.

```
> #
# This routine randomly choose an elliptic "curve" modulo n,
# where gcd(n,6)=1. The coordinates x,y are chosen
# randomly, the parameter a too, and b is calculated.
# The list [x,y,a,b] is given back or a divisor d of n.
#

ellrand:=proc(n) local x,y,a,b,r,d,f;
r:=rand(n); d:=n;
while d=n do
  x:=r(n); y:=r(n); a:=r(n); b:=y^2-x^3-a*x mod n;
  d:=4*a^3+27*b^2 mod n; gcd(d,n);
od;
if %<n and %>1 then return % fi;
[x,y,a,b]; end;
ellrand:=proc(n)
local x, y, a, b, r, d, f;
r:= rand(n);
d:= n;
while d = ndo
  x:= r(n);
  y:= r(n);
  a:= r(n);
  b:= mod(y^2 - x^3 - a*x, n);
  d:= mod(4*a^3 + 27*b^2, n);
  gcd(d, n)
end do;
if % < n and 1 < % then
  return %
end if;
[x, y, a, b]
end proc
(11.1.1)
```

```
> #
# Addition on an elliptic "curve" modulo n, where gcd(n,6)=1.
# P and Q are the points to add and a,b are the parameters.
# The return value is the sum of the points or a divisor d of
n.
```

```

#
elladd:=proc(P,Q,a,b,n) local l,d;
if P[3]=0 then return Q fi;
if Q[3]=0 then return P fi;
if P=Q then return elldou(P,a,b,n) fi;
if P[1]=Q[1] then return [0,1,0] fi;
d:=igcdex(P[1]-Q[1],n,'1');
if 1<d and d<n then return d fi;
l:=(P[2]-Q[2])*l mod n;
l^2-P[1]-Q[1] mod n;
[%,l*(P[1]-%)-P[2] mod n,1];
end;
elladd:=proc(P,Q,a,b,n)
local l,d;
if P[3]=0 then
return Q
end if;
if Q[3]=0 then
return P
end if;
if P=Q then
return elldou(P,a,b,n)
end if;
if P[1]=Q[1] then
return [0,1,0]
end if;
d:=igcdex(P[1]-Q[1],n,'1');
if 1 < d and d < n then
return d
end if;
l:=mod((P[2]-Q[2])*l,n);
mod(l^2-P[1]-Q[1],n);
[%,mod(l*(P[1]-%)-P[2],n),1]
end proc
> #
# Doubling on an elliptic "curve" modulo n, where gcd(n,6)=1.

```

(11.1.2)

```

# P is the point to double and a, b are the parameters.
# The return value is the double of the point P or
# a divisor d of n.
#

```

```

elldou:=proc(P,a,b,n) local l,d;
if P[3]=0 then return P fi;
if P[2]=0 then return [0,1,0] fi;
d:=igcdex(2*P[2],n,'l');
if 1<d and d<n then return d fi;
l:=(3*P[1]^2+a)*l mod n;
l^2-2*P[1] mod n;
[%,l*(P[1]-%) - P[2] mod n,1];
end;

```

```

elldou:=proc(P, a, b, n)

```

(11.1.3)

```

local l, d;
if P[3] = 0 then
    return P
end if;
if P[2] = 0 then
    return [0, 1, 0]
end if;
d := igcdex(2 * P[2], n, 'l');
if 1 < d and d < n then
    return d
end if;
l := mod((3 * P[1]^2 + a) * l, n);
mod(l^2 - 2 * P[1], n);
[%, mod(l * (P[1] - %) - P[2], n), 1]
end proc

```

```

> #
# This program compute k*P, k>=0 or a divisor of n
# on an elliptic "curve" modulo n, where gcd(n,6)=1.
# It use the left-to-right binary method.
#

```

```

ellmul:=proc(P,k,a,b,n) local L,i,Q;
if k=0 then return [0,1,0] fi;
if P[3]=0 then return P fi;
L:=convert(k,base,2);

```



```

Q:=P;
for i from nops(L)-1 to 1 by -1 do
  Q:=elldou(Q,a,b,n);
  if type(Q,integer) then return Q fi;
  if L[i]=1 then
    Q:=elladd(P,Q,a,b,n);
    if type(Q,integer) then return Q fi;
  fi;
od; Q; end;
ellmul:=proc(P, k, a, b, n)
  local L, i, Q;
  if k = 0 then
    return [0, 1, 0]
  end if;
  if P[3] = 0 then
    return P
  end if;
  L:=convert(k, base, 2);
  Q:=P;
  for i from nops(L) - 1 by -1 to 1 do
    Q:=elldou(Q, a, b, n);
    if type(Q, integer) then
      return Q
    end if;
    if L[i] = 1 then
      Q:=elladd(P, Q, a, b, n);
      if type(Q, integer) then
        return Q
      end if
    end if
  end do;
  Q
end proc

```

(11.1.4)

```

> #
# This brute force procedure calculate the number of points
# on an elliptic curve modulo p>3, a prime. The curve

```

```

# parameters are a, b.
#

ellcount:=proc(a,b,p) local x,c;
c:=1;
for x from 0 to p-1 do c:=c+numtheory[jacobi](x^3+a*x+b,p)+1;
od;
c; end;
ellcount:=proc(a, b, p)

```

(11.1.5)

```

local x, c;
c:= 1;
for x from 0 to p - 1 do
    c:= c + numtheory[numtheory:-jacobi](x^3 + a*x + b, p) + 1
end do;
c
end proc

> n:=nextprime(1008); E:=ellrand(n); m:=ellcount(E[3],E[4],n);
n:= 1009
E:= [453, 76, 621, 422]
m:= 1047

```

(11.1.6)

```

> n:=1009; E := [889, 300, 991, 649]; m := 974;
n:= 1009
E:= [889, 300, 991, 649]
m:= 974

```

(11.1.7)

```

> ifactor(974); ellmul([E[1],E[2],1],m,E[3],E[4],n);
elldou([E[1],E[2],1],E[3],E[4],n);
(2) (487)
[0, 1, 0]
[758, 824, 1]

```

(11.1.8)

▼ 11.2. Tétel.

```

> #
# This brute force procedure calculate the high of
# an elliptic curve modulo p>3, a prime. The curve
# parameters are a, b.
#

ellhigh:=proc(a,b,p) local x,y,h,i,P;

```

```

h:=1;
for x from 0 to p-1 do
  y:=msqrt(x^3+a*x+b,p);
  if y=FAIL then next fi;
  P:=[x,y,1];
  for i from 2 do
    i;
    P:=elladd(P,[x,y,1],a,b,p);
    if P=[0,1,0] then
      if i>h then h:=i fi; break;
    fi;
  od; if 3*h>2*p then break fi;
od; h; end;
ellhigh:= proc(a, b, p)
  local x, y, h, i, P;
  h:= 1;
  for x from 0 to p - 1 do
    y:= numtheory.-msqrt(x^3 + a*x + b, p);
    if y = FAIL then
      next
    end if;
    P:= [x, y, 1];
    for i from 2 do
      i;
      P:= elladd(P, [x, y, 1], a, b, p);
      if P = ([0, 1, 0]) then
        if h < i then
          h:= i
        end if;
        break
      end if
    end do;
    if 2 * p < 3 * h then
      break
    end if
  end do;
  h

```

(11.2.1)

```
end proc
```

```
> ellrand(97);  
[58, 15, 0, 83] (11.2.2)
```

```
> ellhigh(5,57,97); ellhigh(95,78,97);  
ellhigh(33,96,97); ellhigh(55,51,97);  
ellhigh(24,13,97); ellhigh(59,39,97);  
ellhigh(49,54,97); ellhigh(45,68,97);  
ellhigh(76,30,97); ellhigh(0,83,97);  
103  
80  
116  
44  
86  
88  
110  
98  
102  
103 (11.2.3)
```

▼ 11.3. A prímtesztek vázlatja.

```
> n:=1009; E:=ellrand(n); m:=ellcount(E[3],E[4],n);  
n:= 1009  
E:= [38, 603, 582, 65]  
m:= 960 (11.3.1)
```

```
> E:=[889,300,991,649]; m:=974;  
E:= [889, 300, 991, 649]  
m:= 974 (11.3.2)
```

```
> ifactor(m); ellmul([E[1],E[2],1],m,E[3],E[4],n);  
elldou([E[1],E[2],1],E[3],E[4],n);  
(2) (487)  
[0, 1, 0]  
[758, 824, 1] (11.3.3)
```

```
> n1:=487; E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);  
n1:= 487  
E1:= [154, 481, 71, 52] (11.3.4)
```

```
m1:=465 (11.3.4)
```

```
> E1:=[201,473,457,19]; m1:=530;  
E1:= [201, 473, 457, 19]
```

```
m1:=530 (11.3.5)
```

```
> ifactor(m1); ellmul([E1[1],E1[2],1],m1,E1[3],E1[4],n1);  
ellmul([E1[1],E1[2],1],10,E1[3],E1[4],n1);  
(2) (5) (53)
```

```
[0, 1, 0]  
[400, 454, 1] (11.3.6)
```

```
> n2:=53;
```

```
n2:=53 (11.3.7)
```

▼ 11.4. A Goldwasser–Kilian–teszt.

```
> n:=1009; E:=ellrand(n); m:=ellcount(E[3],E[4],n);  
n:=1009
```

```
E:= [627, 750, 129, 78]  
m:=1012 (11.4.1)
```

```
> E:=[889,300,991,649]; m:=974;  
E:= [889, 300, 991, 649]
```

```
m:=974 (11.4.2)
```

```
> ifactor(m); ellmul([E[1],E[2],1],m,E[3],E[4],n);  
elldou([E[1],E[2],1],E[3],E[4],n);  
(2) (487)
```

```
[0, 1, 0]  
[758, 824, 1] (11.4.3)
```

```
> n1:=487; E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);  
n1:=487
```

```
E1:= [72, 340, 296, 91]  
m1:=520 (11.4.4)
```

```
> E1:=[201,473,457,19]; m1:=530;  
E1:= [201, 473, 457, 19]
```

```
m1:=530 (11.4.5)
```

```
> ifactor(m1); ellmul([E1[1],E1[2],1],m1,E1[3],E1[4],n1);  
elldou([E1[1],E1[2],1],E1[3],E1[4],n1);  
(2) (5) (53)
```

$$\begin{aligned} & [0, 1, 0] \\ & [68, 198, 1] \end{aligned} \quad (11.4.6)$$

$$\begin{aligned} > \mathbf{n2:=265; E2:=ellrand(n2); m2:=ellcount(E2[3],E2[4],n2);} \\ & \quad n2:=265 \\ & \quad E2:= [230, 259, 171, 136] \\ & \quad m2:=266 \end{aligned} \quad (11.4.7)$$

$$\begin{aligned} > \mathbf{E2:=[230,259,171,136]; m2:=266;} \\ & \quad E2:= [230, 259, 171, 136] \\ & \quad m2:=266 \end{aligned} \quad (11.4.8)$$

$$\begin{aligned} > \mathbf{ifactor(m2); ellmul([E2[1],E2[2],1],m2,E2[3],E2[4],n2);} \\ & \quad \mathbf{elldou([E2[1],E2[2],1],E2[3],E2[4],n2);} \\ & \quad (2) (7) (19) \\ & \quad [208, 89, 1] \\ & \quad [169, 198, 1] \end{aligned} \quad (11.4.9)$$

$$\begin{aligned} > \mathbf{E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);} \\ & \quad E1:= [268, 146, 265, 300] \\ & \quad m1:=451 \end{aligned} \quad (11.4.10)$$

$$\begin{aligned} > \mathbf{E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);} \\ & \quad E1:= [142, 191, 8, 66] \\ & \quad m1:=472 \end{aligned} \quad (11.4.11)$$

$$\begin{aligned} > \mathbf{E1:=[142,191,8,66]; m1:=472;} \\ & \quad E1:= [142, 191, 8, 66] \\ & \quad m1:=472 \end{aligned} \quad (11.4.12)$$

$$\begin{aligned} > \mathbf{ifactor(m1); ellmul([E1[1],E1[2],1],m1,E1[3],E1[4],n1);} \\ & \quad \mathbf{elldou([E1[1],E1[2],1],E1[3],E1[4],n1);} \\ & \quad (2)^3 (59) \\ & \quad [0, 1, 0] \\ & \quad [176, 105, 1] \end{aligned} \quad (11.4.13)$$

$$\begin{aligned} > \mathbf{E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);} \\ & \quad E1:= [249, 267, 252, 400] \\ & \quad m1:=475 \end{aligned} \quad (11.4.14)$$

$$\begin{aligned} > \mathbf{E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);} \\ & \quad E1:= [307, 255, 408, 332] \\ & \quad m1:=522 \end{aligned} \quad (11.4.15)$$

$$> \mathbf{E1 := [307, 255, 408, 332]; m1 := 522;}$$

$$E1 := [307, 255, 408, 332]$$

$$m1 := 522 \quad (11.4.16)$$

> **ifactor(m1); ellmul([E1[1],E1[2],1],m1,E1[3],E1[4],n1);
elldou([E1[1],E1[2],1],E1[3],E1[4],n1);**

$$(2) (3)^2 (29)$$

$$[0, 1, 0]$$

$$[161, 49, 1] \quad (11.4.17)$$

> **E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);**
 $E1 := [246, 327, 360, 30]$
 $m1 := 468 \quad (11.4.18)$

> **E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);**
 $E1 := [112, 89, 273, 301]$
 $m1 := 520 \quad (11.4.19)$

> **E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);**
 $E1 := [170, 371, 182, 391]$
 $m1 := 506 \quad (11.4.20)$

> **E1 := [170, 371, 182, 391]; m1 := 506;**
 $E1 := [170, 371, 182, 391]$
 $m1 := 506 \quad (11.4.21)$

> **ifactor(m1); ellmul([E1[1],E1[2],1],m1,E1[3],E1[4],n1);
elldou([E1[1],E1[2],1],E1[3],E1[4],n1);**

$$(2) (11) (23)$$

$$[0, 1, 0]$$

$$[72, 50, 1] \quad (11.4.22)$$

> **E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);**
 $E1 := [108, 39, 16, 437]$
 $m1 := 474 \quad (11.4.23)$

> **E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);**
 $E1 := [453, 307, 336, 338]$
 $m1 := 465 \quad (11.4.24)$

> **E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);**
 $E1 := [214, 417, 212, 484]$
 $m1 := 528 \quad (11.4.25)$

> **E1:=ellrand(n1); m1:=ellcount(E1[3],E1[4],n1);**
 $E1 := [136, 360, 195, 222]$
 $(11.4.26)$

$m1 := 492$ (11.4.26)

> $E1 := \text{ellrand}(n1); m1 := \text{ellcount}(E1[3], E1[4], n1);$
 $E1 := [338, 124, 446, 259]$

$m1 := 469$ (11.4.27)

> $E1 := \text{ellrand}(n1); m1 := \text{ellcount}(E1[3], E1[4], n1);$
 $E1 := [355, 453, 291, 475]$

$m1 := 525$ (11.4.28)

> $E1 := \text{ellrand}(n1); m1 := \text{ellcount}(E1[3], E1[4], n1);$
 $E1 := [163, 265, 295, 369]$

$m1 := 518$ (11.4.29)

> $E1 := [163, 265, 295, 369]; m1 := 518;$
 $E1 := [163, 265, 295, 369]$

$m1 := 518$ (11.4.30)

> $\text{ifactor}(m1); \text{ellmul}([E1[1], E1[2], 1], m1, E1[3], E1[4], n1);$
 $\text{elldou}([E1[1], E1[2], 1], E1[3], E1[4], n1);$
 $(2) (7) (37)$

$[0, 1, 0]$

$[383, 148, 1]$ (11.4.31)

> $E1 := \text{ellrand}(n1); m1 := \text{ellcount}(E1[3], E1[4], n1);$
 $E1 := [321, 86, 143, 358]$

$m1 := 480$ (11.4.32)

> $E1 := \text{ellrand}(n1); m1 := \text{ellcount}(E1[3], E1[4], n1);$
 $E1 := [353, 426, 308, 26]$

$m1 := 468$ (11.4.33)

> $E1 := \text{ellrand}(n1); m1 := \text{ellcount}(E1[3], E1[4], n1);$
 $E1 := [60, 46, 283, 461]$

$m1 := 520$ (11.4.34)

> $E1 := \text{ellrand}(n1); m1 := \text{ellcount}(E1[3], E1[4], n1);$
 $E1 := [458, 2, 388, 94]$

$m1 := 503$ (11.4.35)

> $E1 := \text{ellrand}(n1); m1 := \text{ellcount}(E1[3], E1[4], n1);$
 $E1 := [365, 230, 106, 399]$

$m1 := 502$ (11.4.36)

> $E1 := [365, 230, 106, 399]; m1 := 502;$
 $E1 := [365, 230, 106, 399]$

(11.4.37)

$m1 := 502$ (11.4.37)

> ifactor(m1); ellmul([E1[1], E1[2], 1], m1, E1[3], E1[4], n1);
elldou([E1[1], E1[2], 1], E1[3], E1[4], n1);

(2) (251)

[0, 1, 0]

[132, 185, 1]

(11.4.38)

> n2:=251; E2:=ellrand(n2); m2:=ellcount(E2[3], E2[4], n2);

n2:= 251

E2:= [10, 164, 111, 188]

m2:= 261

(11.4.39)

> E2:=ellrand(n2); m2:=ellcount(E2[3], E2[4], n2);

E2:= [74, 3, 218, 82]

m2:= 264

(11.4.40)

> E2:=ellrand(n2); m2:=ellcount(E2[3], E2[4], n2);

E2:= [149, 39, 57, 36]

m2:= 252

(11.4.41)

> E2:=ellrand(n2); m2:=ellcount(E2[3], E2[4], n2);

E2:= [220, 225, 106, 119]

m2:= 257

(11.4.42)

> E2:=ellrand(n2); m2:=ellcount(E2[3], E2[4], n2);

E2:= [10, 29, 133, 17]

m2:= 275

(11.4.43)

> E2:=ellrand(n2); m2:=ellcount(E2[3], E2[4], n2);

E2:= [31, 39, 115, 42]

m2:= 232

(11.4.44)

> E2:=ellrand(n2); m2:=ellcount(E2[3], E2[4], n2);

E2:= [151, 79, 95, 195]

m2:= 237

(11.4.45)

> E2:=ellrand(n2); m2:=ellcount(E2[3], E2[4], n2);

E2:= [10, 22, 168, 63]

m2:= 234

(11.4.46)

> E2 := [10, 22, 168, 63]; m2 := 234;

E2:= [10, 22, 168, 63]

m2:= 234

(11.4.47)

> ifactor(m2); ellmul([E2[1], E2[2], 1], m2, E2[3], E2[4], n2);

$$\begin{aligned}
& \text{ell}(\text{dou}([E2[1], E2[2], 1], E2[3], E2[4], n2); \\
& \quad (2) (3)^2 (13) \\
& \quad [0, 1, 0] \\
& \quad [201, 137, 1]
\end{aligned} \tag{11.4.48}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [179, 127, 57, 164] \\
m2 := 281
\end{aligned} \tag{11.4.49}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [66, 241, 120, 111] \\
m2 := 229
\end{aligned} \tag{11.4.50}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [208, 64, 196, 165] \\
m2 := 264
\end{aligned} \tag{11.4.51}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [129, 35, 188, 179] \\
m2 := 240
\end{aligned} \tag{11.4.52}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [83, 223, 221, 2] \\
m2 := 264
\end{aligned} \tag{11.4.53}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [158, 208, 30, 23] \\
m2 := 265
\end{aligned} \tag{11.4.54}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [53, 66, 240, 136] \\
m2 := 234
\end{aligned} \tag{11.4.55}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [186, 230, 65, 179] \\
m2 := 267
\end{aligned} \tag{11.4.56}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [139, 176, 100, 89] \\
m2 := 256
\end{aligned} \tag{11.4.57}$$

$$\begin{aligned}
> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2); \\
E2 := [89, 34, 14, 0] \\
m2 := 252
\end{aligned} \tag{11.4.58}$$

$$> E2 := \text{ellrand}(n2); \quad m2 := \text{ellcount}(E2[3], E2[4], n2);$$

$E2 := [153, 99, 117, 125]$
 $m2 := 246$ (11.4.59)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [151, 135, 190, 93]$
 $m2 := 240$ (11.4.60)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [77, 150, 237, 20]$
 $m2 := 276$ (11.4.61)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [72, 149, 240, 141]$
 $m2 := 280$ (11.4.62)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [159, 225, 8, 243]$
 $m2 := 249$ (11.4.63)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [52, 2, 97, 183]$
 $m2 := 272$ (11.4.64)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [68, 124, 130, 81]$
 $m2 := 270$ (11.4.65)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [72, 215, 164, 19]$
 $m2 := 259$ (11.4.66)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [187, 109, 93, 112]$
 $m2 := 241$ (11.4.67)

> $E2 := \text{ellrand}(n2); m2 := \text{ellcount}(E2[3], E2[4], n2);$
 $E2 := [179, 15, 156, 173]$
 $m2 := 262$ (11.4.68)

> $E2 := [179, 15, 156, 173]; m2 := 262;$
 $E2 := [179, 15, 156, 173]$
 $m2 := 262$ (11.4.69)

> $\text{ifactor}(m2); \text{ellmul}([E2[1], E2[2], 1], m2, E2[3], E2[4], n2);$
 $\text{elldou}([E2[1], E2[2], 1], E2[3], E2[4], n2);$
(2) (131)

```
[0, 1, 0]
```

```
[219, 125, 1] (11.4.70)
```

```
> n3:=131; E3:=ellrand(n3); m3:=ellcount(E3[3],E3[4],n3);
```

```
n3:= 131
```

```
E3:= [50, 130, 70, 12]
```

```
m3:= 127 (11.4.71)
```

```
> E3:=ellrand(n3); m3:=ellcount(E3[3],E3[4],n3);
```

```
E3:= [104, 73, 48, 108]
```

```
m3:= 113 (11.4.72)
```

```
> E3:=ellrand(n3); m3:=ellcount(E3[3],E3[4],n3);
```

```
E3:= [32, 128, 104, 69]
```

```
m3:= 146 (11.4.73)
```

```
> E3 := [32, 128, 104, 69]; m3 := 146;
```

```
E3:= [32, 128, 104, 69]
```

```
m3:= 146 (11.4.74)
```

```
> ifactor(m3); ellmul([E3[1],E3[2],1],m3,E3[3],E3[4],n3);
```

```
elldou([E3[1],E3[2],1],E3[3],E3[4],n3);
```

```
(2) (73)
```

```
[0, 1, 0]
```

```
[110, 26, 1] (11.4.75)
```

```
> n4:=73;
```

▼ 11.5. Atkin tesztje.

```
> interface(echo=2);
```

```
1
```

```
(11.5.1)
```

```
> ;
```

```
> #
```

```
# For a given negative discriminant -D
```

```
# we calculate the ideal class number.
```

```
#
```

```
idealclassnumber:=proc(D) local a,b,c,n,h; h:=0;
```

```
for a while 3*a^2<=D do
```

```
  b:=D mod 2; n:=(b^2+D)/4;
```

```
  while b<=a do
```

```
    if n mod a=0 then c:=n/a;
```

```
    if c>=a and igcd(a,b,c)=1 then
```

```
      if a=c or b=0 or b=a then h:=h+1 else h:=h+2 fi;
```

```

        fi;
        fi; n:=n+b+1; b:=b+2;
    od;
od; h; end;
idealclassnumber:=proc(D)
local a, b, c, n, h;
h:=0;
for a while 3*a^2 <= D do
    b:=mod(D, 2);
    n:=1/4*b^2+1/4*D;
    while b <= a do
        if mod(n, a) = 0 then
            c:=n/a;
            if a <= c and igcd(a, b, c) = 1 then
                if a = c or b = 0 or b = a then
                    h:=h+1
                else
                    h:=h+2
                end if
            end if
        end if;
        n:=n+b+1;
        b:=b+2
    end do
end do;
h
end proc
> #
# For a given negative discriminant -D
# this calculates the ideal class number.
#
# This versions is slower then the other!
#
idealclassnumber2:=proc(D) local a,b,c,n,h,X,x,DD; h:=1;
if type(D,even) then
    DD:=D/4;

```

```

for a from 2 while 3*a^2<=D do
  X:=Roots(x^2+DD) mod a;
  for b in X do b:=2*b[1];
    if b>a then b:=b-2*a; fi;
    c:=(b^2+D)/4/a;
    if c>=a and igcd(a,b,c)=1 then
      if a=c and b<0 then next fi;
      h:=h+1;
    fi;
  od;
od;
else
DD:=(D+1)/4;
for a from 2 while 3*a^2<=D do
  X:=Roots(x^2+x+DD) mod a;
  for b in X do
    b:=2*b[1]+1;
    if b>a then b:=b-2*a; fi;
    c:=(b^2+D)/4/a;
    if c>=a and igcd(a,b,c)=1 then
      if a=c and b<0 then next fi;
      h:=h+1;
    fi;
  od;
od;
fi; h; end;
idealclassnumber2:= proc(D)
local a, b, c, n, h, X, x, DD;
h:= 1;
if type(D, even) then
  DD:= 1 / 4 * D;
  for a from 2 while 3 * a^2 <= D do
    X:= mod(Roots(x^2 + DD), a);
    for b in X do
      b:= 2 * b[1];
      if a < b then
        b:= b - 2 * a
      end if;
      c:= 1 / 4 * (b^2 + D) / a;
      if a <= c and igcd(a, b, c) = 1 then
        if a = c and b < 0 then

```

```

        next
        end if;
        h:= h+1
    end if
end do
end do
else
    DD:= 1 / 4 * D + 1 / 4;
    for a from 2 while 3 * a^2 <= D do
        X:= mod(Roots(x^2 + x + DD), a);
        for b in X do
            b:= 2 * b[1] + 1;
            if a < b then
                b:= b - 2 * a
            end if;
            c:= 1 / 4 * (b^2 + D) / a;
            if a <= c and igcd(a, b, c) = 1 then
                if a = c and b < 0 then
                    next
                end if;
                h:= h+1
            end if
        end do
    end do
end if;
h
end proc

```

> #
For a given negative discriminant -D
this calculates the ideal class number.

This versions is faster as others for
large discriminants

NOT YET FINISHED!!!!

```

#
idealclassnumbershank:=proc(D,f,b) local P,p,Q,d,DD;
DD:=2^16; # have to adjust for an appropriate limit
if D<DD then return(idealclassnumber(D)) fi;
P:=ceil(D^(1/4.));
Q:=1.*sqrt(D)/Pi/(1.-numtheory[jacobi](-D,2)/2.);
Q:=Q/(1.-numtheory[jacobi](-D,3)/3.);
p:=5; d:=2;
while p<=P do
  if isprime(p) then Q:=Q/(1.-numtheory[jacobi](-D,p)/p); fi;
  p:=p+d; d:=6-d;
od; end;
idealclassnumbershank:=proc(D, f, b)
local P, p, Q, d, DD;
DD:= 65536;
if D < DD then
  return idealclassnumber(D)
end if;
P:= ceil(D^(1 / 4.));
Q:= 1.*sqrt(D) / (pi*(1. - numtheory[jacobi](-D, 2) / 2.));
Q:= Q / (1 - numtheory[jacobi](-D, 3) / 3.);
p:= 5;
d:= 2;
while p <= P do
  if isprime(p) then
    Q:= Q / (1 - numtheory[jacobi](-D, p) / p)
  end if;
  p:= p + d;
  d:= 6 - d
end do
end proc
> #
# This procedure gives back the next fundamental
# discriminant larger then d. A discriminant have
# to be larger then 6, congruent to 3 modulo 4 or congruent
# to 4 or 8 modulo 16 and a square of an odd prime
# may not divide it. The result is in form
# [d,h,q1,q2,...], where d is the discriminant,

```



```

# h is the ideal class number in  $Q(\sqrt{-d})$ ,
# and the q's are odd primes except maybe q1, which may 4 or
# 8.
#

```

```

nextdiscriminant:=proc(d) local dd,dp,L,n,p,i;
dd:=max(d+1,7);
while true do
  if dd mod 16=4 or dd mod 16=8 or dd mod 4=3 then
    L:=[]; n:=dd;
    if type(n,even) then
      for i from 0 while type(n,even) do n:=n/2; od;
      L:=[2^i];
    fi;
    if n>=9 then
      if n mod 3=0 then n:=n/3;
        if n mod 3=0 then dd:=dd+1; next; fi;
        L:=[op(L),3];
      fi;
    fi;
    dp:=2; p:=5;
    while n>=p^2 do
      if n mod p=0 then n:=n/p;
        if n mod p=0 then L:=false; break; fi;
        L:=[op(L),p];
      fi;
      p:=p+dp; dp:=6-dp;
    od;
    if L=false then dd:=dd+1; next; fi;
    L:=[dd,idealclassnumber(dd),op(L)];
    if n>1 then L:=[op(L),n] fi;
    break;
  else
    dd:=dd+1;
  fi;
od; L; end;

```

nextdiscriminant:= proc(d)

(11.5.5)

local dd, dp, L, n, p, i;

dd:= max(d + 1, 7);

do

if mod(dd, 16) = 4 or mod(dd, 16) = 8 or mod(dd, 4) = 3 then

L:= [];

n:= dd;

if type(n, even) then

```

for  $i$  from 0 while  $\text{type}(n, \text{even})$  do
     $n := 1 / 2 * n$ 
end do;
 $L := [2^i]$ 
end if;
if  $9 \leq n$  then
    if  $\text{mod}(n, 3) = 0$  then
         $n := 1 / 3 * n;$ 
        if  $\text{mod}(n, 3) = 0$  then
             $dd := dd + 1;$ 
        next
        end if;
         $L := [\text{op}(L), 3]$ 
    end if
end if;
 $dp := 2;$ 
 $p := 5;$ 
while  $p^2 \leq n$  do
    if  $\text{mod}(n, p) = 0$  then
         $n := n / p;$ 
        if  $\text{mod}(n, p) = 0$  then
             $L := \text{false};$ 
            break
        end if;
         $L := [\text{op}(L), p]$ 
    end if;
     $p := p + dp;$ 
     $dp := 6 - dp$ 
end do;
if  $L = \text{false}$  then
     $dd := dd + 1;$ 
next

```

```

        end if;
        L:= [dd, idealclassnumber(dd), op(L)];
        if 1 < n then
            L:= [op(L), n]
        end if;
        break
    else
        dd:= dd + 1
    end if
end do;
L
end proc
> #
# This procedure write to a file all the fundamental
# discriminants
# in the interval II. The ideal class number h also
# calculated,
# and written. First h is written on two bytes, then the
# number
# of the factors on one byte. The last factor is written on 4
# bytes, the second and third last factors are written on 2
# bytes,
# all other factors are written on one byte.
#

discriminants:=proc(II,filename) local x,L,B,D,id;
x:=op(1,II)-1;
if x<7 then
    id:=fopen(filename,WRITE,BINARY)
else
    id:=fopen(filename,APPEND,BINARY)
fi;
B:=op(2,II); x:=nextdiscriminant(x); D:=x[1];
while D<=B do
    writediscriminant(x,id); x:=nextdiscriminant(D); D:=x[1];
od;
fclose(id);
end;
discriminants:= proc(II, filename)
    local x, L, B, D, id;
    x:= op(1, II) - 1;

```

(11.5.6)

```

if  $x < 7$  then
     $id := fopen(filename, WRITE, BINARY)$ 
else
     $id := fopen(filename, APPEND, BINARY)$ 
end if;
 $B := op(2, I);$ 
 $x := nextdiscriminant(x);$ 
 $D := x[1];$ 
while  $D \leq B$  do
     $writediscriminant(x, id);$ 
     $x := nextdiscriminant(D);$ 
     $D := x[1]$ 
end do;
 $fclose(id)$ 
end proc

```

```

> #
# This procedure write to a file a discriminant having size
# up to ds bytes; if ds is not given, ds=4 is the default
# value.
# First  $h(-d)$  is written on hs bytes; if hs is not given,
# we presuppose  $hs = \lceil ds/2 + 1/2 \rceil$ , which is large enough
# to contain the ideal class number  $h(-d) < \sqrt{d} \ln(\sqrt{d})$ 
# /Pi
# whenever  $ds \leq 18$ . Then the number of the factors is written
# on one byte, and the factors starting with 4 or 8, if
# there is such a factor and continuing with odd factors in
# increasing order. The last factor is written on ds bytes,
# the second last factor on  $\lceil ds/2 \rceil$  bytes, the third last
# factors on  $\lceil ds/3 \rceil$  bytes, etc. If ps is also given, all
# factors are written on at most ps bytes. Discriminants
# which
# cannot included in this format cause error.
#

writediscriminant:=proc(L,id) local i,j,k,l,B,ds,hs,ps;
if nargs<2 then ERROR(`not enough args`) fi;
if nargs=2 then ds:=4 else ds:=args[3] fi;
if nargs<=3 then hs:=ceil((ds/2+1/2)) else hs:=args[4] fi;
if nargs<=4 then ps:=ds else ps:=args[5] fi;
B:=convert(L[2],base,256);
if nops(B)>hs then ERROR(`too large ideal class number`,L[2])

```

```

fi;
if nops(B)<hs then B:=[op(B),0$j=nops(B)+1..hs] fi;
writebytes(id,B);
i:=3; k:=nops(L)-2;
writebytes(id,[k]);
while k>=ds do writebytes(id,[L[i]]); i:=i+1; k:=k-1; od;
while k>0 do l:=ceil(ds/k); if l>ps then l:=ps fi;
  B:=convert(L[i],base,256);
  if nops(B)>l then ERROR(`too large prime factor`,L[i]) fi;
  if nops(B)<l then B:=[op(B),0$j=nops(B)+1..l] fi;
  writebytes(id,B); i:=i+1; k:=k-1;
od; end;

```

writediscriminant:= proc(L, id)

(11.5.7)

local i, j, k, l, B, ds, hs, ps;

if nargs < 2 then

ERROR(not enough args)

end if;

if nargs = 2 then

ds:= 4

else

ds:= args[3]

end if;

if nargs <= 3 then

hs:= ceil(1 / 2 * ds + 1 / 2)

else

hs:= args[4]

end if;

if nargs <= 4 then

ps:= ds

else

ps:= args[5]

end if;

B:= convert(L[2], base, 256);

if hs < nops(B) then

ERROR(too large ideal class number, L[2])

end if;

```

if  $nops(B) < hs$  then
     $B := [op(B), \$(0, j = nops(B) + 1 ..hs)]$ 
end if;
writebytes( $id, B$ );
 $i := 3$ ;
 $k := nops(L) - 2$ ;
writebytes( $id, [k]$ );
while  $ds \leq k$  do
    writebytes( $id, [L[i]]$ );
     $i := i + 1$ ;
     $k := k - 1$ 
end do;
while  $0 < k$  do
     $l := \text{ceil}(ds / k)$ ;
    if  $ps < l$  then
         $l := ps$ 
    end if;
     $B := \text{convert}(L[i], \text{base}, 256)$ ;
    if  $l < nops(B)$  then
        ERROR(too large prime factor, L[i])
    end if;
    if  $nops(B) < l$  then
         $B := [op(B), \$(0, j = nops(B) + 1 ..l)]$ 
    end if;
    writebytes( $id, B$ );
     $i := i + 1$ ;
     $k := k - 1$ 
end do
end proc

```

```

> #
# This procedure read from a file a discriminant having size
# up to ds bytes; if ds is not given, ds=4 is the default
# value.
# First h(-d) is read on hs bytes; if hs is not given,

```

```

# we presuppose  $hs = \text{ceil}(ds/2 + 1/2)$ , which is large enough
# to contain the ideal class number  $h(-d) < \sqrt{d} \ln(\sqrt{d}) / \pi$ 
# whenever  $ds < 19$ . Then the number of the factors on one byte
# is read, and the factors starting with 4 or 8, if there is
# such a factor and continuing with odd factors in increasing
# order. The last factor is supposed to be written on  $ds$ 
# bytes,
# the second last factor on  $\text{ceil}(ds/2)$  bytes, the third last
# factor on  $\text{ceil}(ds/3)$  bytes, etc. If  $ps$  is also given, all
# factors are supposed to be written on at most  $ps$  bytes.
#

```

```

readdiscriminant:=proc(id)
local i,j,k,B,l,d,L,ds,hs,ps;
if nargs<1 then ERROR(`not enough args`) fi;
if nargs=1 then ds:=4 else ds:=args[2] fi;
if nargs<=2 then hs:=ceil((ds/2+1/2)) else hs:=args[3] fi;
if nargs<=3 then ps:=ds else ps:=args[4] fi;
B:=readbytes(id,hs); if B=0 then RETURN(NULL) fi;
L:=[convert([B[j]*256^(j-1)$j=1..nops(B)],`+`)];
B:=readbytes(id,1);k:=B[1]; d:=1;
while k>=ds do
  B:=readbytes(id,1); L:=[op(L),B[1]];
  d:=d*B[1]; k:=k-1;
od;
while k>0 do l:=ceil(ds/k); if l>ps then l:=ps fi;
  B:=readbytes(id,l);
  B:=convert([B[j]*256^(j-1)$j=1..nops(B)],`+`);
  d:=d*B; L:=[op(L),B]; k:=k-1;
od; [d,op(L)]; end;

```

readdiscriminant:= **proc**(*id*)

(11.5.8)

```

local i, j, k, B, l, d, L, ds, hs, ps;
if nargs < 1 then
  ERROR(not enough args)
end if;
if nargs = 1 then
  ds := 4
else
  ds := args[2]
end if;
if nargs <= 2 then
  hs := ceil(1 / 2 * ds + 1 / 2)

```

```

else
     $hs := args[3]$ 
end if;
if  $nargs \leq 3$  then
     $ps := ds$ 
else
     $ps := args[4]$ 
end if;
 $B := readbytes(id, hs)$ ;
if  $B = 0$  then
    RETURN(NULL)
end if;
 $L := [convert([\$(B[j]*256^{(j-1)}, j = 1..nops(B)]), +)]$ ;
 $B := readbytes(id, 1)$ ;
 $k := B[1]$ ;
 $d := 1$ ;
while  $ds \leq k$  do
     $B := readbytes(id, 1)$ ;
     $L := [op(L), B[1]]$ ;
     $d := d*B[1]$ ;
     $k := k - 1$ 
end do;
while  $0 < k$  do
     $l := \text{ceil}(ds/k)$ ;
    if  $ps < l$  then
         $l := ps$ 
    end if;
     $B := readbytes(id, l)$ ;
     $B := convert([\$(B[j]*256^{(j-1)}, j = 1..nops(B)]), +)$ ;
     $d := d*B$ ;
     $L := [op(L), B]$ ;
     $k := k - 1$ 
end while;

```



```

end do;
[d, op(L)]
end proc

> #
# This procedure sieve out discriminants from a file for
# which
# the function f gives value true. The result is [n,e] where
# n
# is the number of good discriminants and e is the sum of 1/h
# (D)'s.
#

sievediscriminants:=proc(filename, f) local e, n, x, id;
id:=fopen(filename, READ, BINARY); e:=0.0; n:=0;
do
  x:=readdiscriminant(id);
  if x=NULL then break fi;
  if f(x) then n:=n+1; e:=e+1/x[2]; fi;
od; fclose(id); [n, e]; end;
sievediscriminants:=proc(filename, f)
local e, n, x, id;
id:=fopen(filename, READ, BINARY);
e:=0.;
n:=0;
do
  x:=readdiscriminant(id);
  if x=NULL then
    break
  end if;
  if f(x) then
    n:=n+1;
    e:=e+1/x[2]
  end if
end do;
fclose(id);
[n, e]
end proc

> badness:=x->x[2]/2^(nops(x)-3);

```

(11.5.9)

$$\text{badness} := x \rightarrow \frac{x_2}{2^{\text{nops}(x)} - 3} \quad (11.5.10)$$

2. Find good discriminants

```

> #
# This procedure calculates square root of the number a
# modulo a "probable prime" p. If p is not a prime, an error
# message may be caused.
#

modsqrt:=proc(a::integer,p::posint) local k,s,r,n,ra,j,i,z,c,
x,t,b;
if p=1 or type(p,even) then ERROR(`not a prime in modsqrt`,p)
fi;
if numtheory[jacobi](a,p)<>1 then ERROR(`not a square`,a) fi;
k:=p-1; s:=0;
while type(k,even) do s:=s+1; k:=k/2; od;
k:=(k-1)/2; r:=a^(k+1) mod p;n:=a^(2*k+1) mod p;
ra:=rand(p); j:=1;
for i to 100 while j<>-1 do
z:=ra(); j:=numtheory[jacobi](z,p);
od;
if i=100 then ERROR(`probably not a prime in modsqrt`,p) fi;
c:=z^(2*k+1) mod p;
while n<>1 do x:=n;
for t from s to 0 by -1 while x<>1 do x:=x^2 mod p; od;
if t=0 then ERROR(`not a prime in modsqrt`,p) fi;
b:=c^(2^(t-1)) mod p; r:=r*b mod p; c:=b^2 mod p;
n:=n*c mod p; s:=s-t;
od; r; end;

```

modsqrt := proc(*a*::integer, *p*::posint) (11.5.11)

```

local k, s, r, n, ra, j, i, z, c, x, t, b;
if p = 1 or type(p, even) then
    ERROR(not a prime in modsqrt, p)
end if;
if numtheory[jacobi](a, p) <> 1 then
    ERROR(not a square, a)
end if;
k := p - 1;
s := 0;
while type(k, even) do

```

```

    s := s + 1;
    k := 1 / 2 * k
end do;
k := 1 / 2 * k - 1 / 2;
r := mod(a &^ (k + 1), p);
n := mod(a &^ (2 * k + 1), p);
ra := rand(p);
j := 1;
for i to 100 while j <> -1 do
    z := ra();
    j := numtheory[jacobi](z, p)
end do;
if i = 100 then
    ERROR(probably not a prime in modsqrt, p)
end if;
c := mod(z &^ (2 * k + 1), p);
while n <> 1 do
    x := n;
    for t from s by -1 to 0 while x <> 1 do
        x := mod(x^2, p)
    end do;
    if t = 0 then
        ERROR(not a prime in modsqrt, p)
    end if;
    b := mod(c &^ (2^(t - 1)), p);
    r := mod(r * b, p);
    c := mod(b^2, p);
    n := mod(n * c, p);
    s := s - t
end do;
r
end proc

```

```

> #
# This procedure put all the primes p below P to the
# table goodprimes, for which (n|p)=1, where
# n is a given probable prime. Differences between
# odd primes are read from a file.
#

goodprimes:=proc(filename,P,n,goodprimes) local p,id,d;
if numtheory[jacobi](-1,n)=1 then goodprimes[-1]:=-1
  else goodprimes[-1]:=0; fi;
if numtheory[jacobi](-2,n)=1 then goodprimes[-2]:=-2
  else goodprimes[-2]:=0; fi;
if numtheory[jacobi](2,n)=1 then goodprimes[2]:=2
  else goodprimes[2]:=0 fi;
id:=fopen(filename,READ,BINARY); p:=3;
while p<P do
  if numtheory[jacobi](n,p)=1 then
    goodprimes[p]:=1;
  else
    goodprimes[p]:=0;
  fi;
  d:=readdiff(id);
  if d=NULL then ERROR(`not enough prime differences`) fi;
  p:=p+d;
od; fclose(id); end;
goodprimes:=proc(filename,P,n,goodprimes)
  local p, id, d;
  if numtheory[jacobi](-1,n) = 1 then
    goodprimes[-1] := -1
  else
    goodprimes[-1] := 0
  end if;
  if numtheory[jacobi](-2,n) = 1 then
    goodprimes[-2] := -2
  else
    goodprimes[-2] := 0
  end if;
  if numtheory[jacobi](2,n) = 1 then
    goodprimes[2] := 2
  else
    goodprimes[2] := 0

```

```

end if;
id:= fopen(filename, READ, BINARY);
p:= 3;
while p < P do
    if numtheory[jacobi](n, p) = 1 then
        goodprimes[p]:= 1
    else
        goodprimes[p]:= 0
    end if;
    d:= readdiff(id);
    if d = NULL then
        ERROR(not enough prime differences)
    end if;
    p:= p + d
end do;
fclose(id)
end proc

> #
# This procedure put the square root of p or -p
# for all the odd primes p below P to the
# table roots, for which (n|p)=1, where
# n is a given probable prime. Differences between
# odd primes are read from a file. Good primes came
# from table goodprimes.
#

goodprimeroots:=proc(filename,P,n,goodprimes,roots) local p,
id,d;
if goodprimes[-1]<>0 then roots[-1]:=modsqrt(n-1,n) fi;
if goodprimes[-2]<>0 then roots[-2]:=modsqrt(n-2,n) fi;
if goodprimes[2]<>0 then roots[2]:=modsqrt(2,n) fi;
id:=fopen(filename,READ,BINARY); p:=3;
if n mod 4=1 then
    while p<P do
        if goodprimes[p]<>0 then roots[p]:=modsqrt(p,n); fi;
        d:=readdiff(id);
        if d=NULL then ERROR(`not enough prime differences`) fi;
        p:=p+d;
    od;
else

```

```

while p<P do
  if goodprimes[p]<>0 then
    if p mod 4=1 then
      roots[p]:=modsqrt(p,n);
    else
      roots[p]:=modsqrt(n-p,n);
    fi;
  fi;
  d:=readdiff(id);
  if d=NULL then ERROR(`not enough prime differences`) fi;
  p:=p+d;
od;
fi; fclose(id); end;
goodprimeroots:= proc(filename, P, n, goodprimes, roots)
  local p, id, d;
  if goodprimes[-1]<>0 then
    roots[-1]:= modsqrt(n-1, n)
  end if;
  if goodprimes[-2]<>0 then
    roots[-2]:= modsqrt(n-2, n)
  end if;
  if goodprimes[2]<>0 then
    roots[2]:= modsqrt(2, n)
  end if;
  id:= fopen(filename, READ, BINARY);
  p:= 3;
  if mod(n, 4) = 1 then
    while p < P do
      if goodprimes[p] <> 0 then
        roots[p] := modsqrt(p, n)
      end if;
      d := readdiff(id);
      if d = NULL then
        ERROR(not enough prime differences)
      end if;
      p := p + d
    end do

```

(11.5.13)

```

else
  while p < P do
    if goodprimes[p] <> 0 then
      if mod(p, 4) = 1 then
        roots[p] := modsqrt(p, n)
      else
        roots[p] := modsqrt(n - p, n)
      end if
    end if;
    d := readdiff(id);
    if d = NULL then
      ERROR(not enough prime differences)
    end if;
    p := p + d
  end do
end if;
fclose(id)
end proc
> #
# This procedure gives the next good discriminant for
# the probably prime number n. The function value have to be
# true for the discriminant. The condition (-d|n)=1 is
# checked.
# Good prime roots came from table roots.
#
nextgooddiscriminantfor:=proc(id,n,f,roots) local x,i,j,s,q,
ff;
do
  x:=readdiscriminant(id);
  if x=NULL then return(NULL) fi;
  if f(x) then
    i:=3; q:=x[3]; s:=1; ff:=true;
    if q=4 or q=8 then i:=i+1 else q:=1 fi;
    for j from i to nops(x) do
      if not type(roots[x[j]],integer) then ff:=false; break;
    fi;
    if x[j] mod 4=3 then s:=-s; fi;
  od;

```

```

    if not ff then next; fi;
    i:=n mod 8;
    if i=1 or i=5 then s:=1 else s:=-s; fi;
    if q=8 and (i=3 or i=5) then s:=-s; fi;
    if s=1 then break; fi;
  fi;
od; x; end;
nextgooddiscriminantfor:= proc(id, n, f, roots)
local x, i, j, s, q, ff;
do
  x:= readdiscriminant(id);
  if x = NULL then
    return NULL
  end if;
  if f(x) then
    i:= 3;
    q:= x[3];
    s:= 1;
    ff:= true;
    if q = 4 or q = 8 then
      i:= i + 1
    else
      q:= 1
    end if;
    for j from i to nops(x) do
      if not type(roots[x[j]], integer) then
        ff:= false;
        break
      end if;
      if mod(x[j], 4) = 3 then
        s:= -s
      end if
    end do;
    if not ff then
      next

```

(11.5.14)


```

        end if;
        i:= mod(n, 8);
        if i = 1 or i = 5 then
            s:= 1
        else
            s:= -s
        end if;
        if q = 8 and (i = 3 or i = 5) then
            s:= -s
        end if;
        if s = 1 then
            break
        end if
    end if
end do;
x
end proc

> #
# This procedure sieve out discriminants from a file for
# which
# the function f gives value true, and conditions  $(-d|n)=1$ ,
#  $(n|p)=1$ 
# are satisfied. The result is [m,e] where m is the number of
# good
# discriminants and e is the expected number of elliptic
# curves.
#

gooddiscriminantsfor:=proc(filename,n,f) local e,m,x,id;
id:=fopen(filename,READ,BINARY);
e:=0.0; m:=0;
do
    x:=nextgooddiscriminantfor(id,n,f);
    if x=NULL then break fi;
    m:=m+1; e:=e+2^(nops(x)-2)/x[2];
od; fclose(id); [m,e]; end;
gooddiscriminantsfor:=proc(filename, n, f)
local e, m, x, id;
id:= fopen(filename, READ, BINARY);

```

(11.5.15)

```

e:= 0;
m:= 0;
do
  x:= nextgooddiscriminantfor(id, n, f);
  if x = NULL then
    break
  end if;
  m:= m + 1;
  e:= e + 2^(nops(x) - 2) / x[2]
end do;
fclose(id);
[m, e]
end proc
> #
# This procedure gives the next good discriminant and its
# square
# root for the probably prime number n. The function value
# have
# to be true for the discriminant. The condition (-d|n)=1 is
# checked. Prime roots came from table roots.
#
nextgooddiscriminantroot:=proc(id,n,f,roots) local r,x,i,j,s,
q,ff;
do
  x:=readdiscriminant(id);
  if x=NULL then return(NULL) fi;
  if f(x) then
    i:=3; q:=x[3]; s:=1; ff:=true;
    if q=4 or q=8 then i:=i+1 else q:=1 fi;
    for j from i to nops(x) do
      if roots[x[j]]=0 then ff:=false; break; fi;
      if x[j] mod 4=3 then s:=-s; fi;
    od;
    if not ff then next; fi;
    j:=n mod 8;
    if j=1 or j=5 then s:=1 else s:=-s; fi;
    if q=8 and (j=3 or j=5) then s:=-s; fi;
    if s=1 then r:=1;
      for j from i to nops(x) do r:=r*roots[x[j]] mod n; od;
      if q=4 then r:=r*2*roots[-1] mod n; fi;
      j:=n mod 8;

```

```

    if q=8 then
      if (j=3 or j=5) then
        r:=r*2*roots[2] mod n
      else
        r:=r*2*roots[-2] mod n
      fi;
    fi;
  break;
fi;
od; [x[1],r] end;
nextgooddiscriminantroot:=proc(id, n, f, roots)
local r, x, i, j, s, q, ff;
do
  x:=readdiscriminant(id);
  if x = NULL then
    return NULL
  end if;
  if f(x) then
    i:= 3;
    q:= x[3];
    s:= 1;
    ff:= true;
    if q = 4 or q = 8 then
      i:= i+ 1
    else
      q:= 1
    end if;
    for j from i to nops(x) do
      if roots[x[j]] = 0 then
        ff:= false;
        break
      end if;
      if mod(x[j], 4) = 3 then
        s:= -s
      end if

```

(11.5.16)

```

end do;
if not ff then
    next
end if;
j := mod(n, 8);
if j = 1 or j = 5 then
    s := 1
else
    s := -s
end if;
if q = 8 and (j = 3 or j = 5) then
    s := -s
end if;
if s = 1 then
    r := 1;
    for j from i to nops(x) do
        r := mod(r* roots[x[j]], n)
    end do;
    if q = 4 then
        r := mod(2*r* roots[-1], n)
    end if;
    j := mod(n, 8);
    if q = 8 then
        if j = 3 or j = 5 then
            r := mod(2*r* roots[2], n)
        else
            r := mod(2*r* roots[-2], n)
        end if
    end if;
    break
end if
end if

```

```

end do;
[x[1], r]
end proc

```

3. Reduced form we look for

```

> DD:=7; a:=1; b:=-DD; c:=(-DD)*(-DD-1)/4;
      DD:= 7
      a:= 1
      b:=-7
      c:= 14

```

(11.5.17)

```

> DD:=8; a:=1; b:=-DD; c:=(-DD)*(-DD-1)/4;
      DD:= 8
      a:= 1
      b:=-8
      c:= 18

```

(11.5.18)

4. Reduction of forms

```

> #
# This procedure calculates for a quadratic form a*x*x+b*x*y+
# c*y*y
# the corresponding reduced quadratic form. The values x, y
# for
# which the reduced form takes the same value as the original
# form for x=1, y=0 also calculated. The new coefficients
# and
# x, y are given back.
#
reduceform:=proc(a,b,c) local aa,bb,cc,x,y,q;
aa:=a; bb:=b; cc:=c;
x:=1; y:=0;
while cc<aa or bb<=-aa or bb>aa or (cc=aa and bb<0) do
  q:=bb+aa-1; q:=(q-(q mod (2*aa)))/2/aa;
  bb:=bb-2*aa*q; cc:=cc-q*(bb+q*aa); x:=x+q*y;
  if cc<aa then q:=x; x:=y; y:=q; q:=aa; aa:=cc; cc:=q; fi;
  if cc=aa and bb<0 then bb:=-bb; x:=-x; fi;
od; [aa,bb,cc,x,y]; end;
reduceform:=proc(a,b,c)
  local aa, bb, cc, x, y, q;

```

(11.5.19)

```

aa:= a;
bb:= b;
cc:= c;
x:= 1;
y:= 0;
while cc < aa or bb <= -aa or aa < bb or cc = aa and bb < 0 do
  q:= bb + aa - 1;
  q:= 1 / 2 * (q - (mod(q, 2 * aa))) / aa;
  bb:= bb - 2 * aa * q;
  cc:= cc - q * (bb + aa * q);
  x:= x + q * y;
  if cc < aa then
    q:= x;
    x:= y;
    y:= q;
    q:= aa;
    aa:= cc;
    cc:= q
  end if;
  if cc = aa and bb < 0 then
    bb:= -bb;
    x:= -x
  end if
end do;
[aa, bb, cc, x, y]
end proc
> DD:=7; a:=1; b:=-DD; c:=(-DD)*(-DD-1)/4; reduceform(a,b,c);
      DD:= 7
      a:= 1
      b:= -7
      c:= 14
      [1, 1, 2, 1, 0]

```

(11.5.20)

```

> DD:=8; a:=1; b:=-DD; c:=(-DD)*(-DD-1)/4; reduceform(a,b,c);
      DD:= 8
      a:= 1
      b:=-8
      c:= 18
      [1, 0, 2, 1, 0]

```

(11.5.21)

5. Find the algebraic integer

```

> #
# This procedure gives the next good discriminant and
# the trace of the corresponding algebraic integer for
# the probably prime number n. The function value
# have to be true for the discriminant.
#
nextgoodorder:=proc(id,n,f,roots) local x,t;
do
  x:=nextgooddiscriminantroot(id,n,f,roots);
  if x=NULL then return(NULL) fi;
  t:=testorder(x[1],n);
  if t<>NULL then break fi;
od; [x[1],t] end;
nextgoodorder:=proc(id, n, f, roots)
local x, t;
do
  x:= nextgooddiscriminantroot(id, n, f, roots);
  if x = NULL then
    return NULL
  end if;
  t:= testorder(x[1], n);
  if t <> NULL then
    break
  end if
end do;
[x[1], t]
end proc
> #

```

(11.5.22)

```

# This procedure count good orders for
# the probably prime number n. The function value
# have to be true for the discriminant.
#

```

```

countgoodorders:=proc(discr,n,f,roots) local x,i,id; i:=0;
id:=fopen(discr,READ,BINARY);
do
  x:=nextgoodorder(id,n,f,roots);
  if x=NULL then break fi; i:=i+2;
od; fclose(discr); i; end;

```

countgoodorders := **proc**(*discr*, *n*, *f*, *roots*) (11.5.23)

```

local x, i, id;
i := 0;
id := fopen(discr, READ, BINARY);
do
  x := nextgoodorder(id, n, f, roots);
  if x = NULL then
    break
  end if;
  i := i + 2;
end do;
fclose(discr);
i
end proc

```

```

> #
# This procedure gives back the list of good orders for
# the probably prime number n. The function value
# have to be true for the discriminant.
#

```

```

goodorders:=proc(discr,n,f) local x,i,id,L; i:=0; L:=[];
id:=fopen(discr,READ,BINARY);
do
  x:=nextgoodorder(id,n,f,roots);
  L:=[op(L),x];
  if x=NULL then break fi; i:=i+2;
od; fclose(discr); L; end;

```

goodorders := **proc**(*discr*, *n*, *f*) (11.5.24)

```

local x, i, id, L;

```



```

i:= 0;
L:= [];
id:= fopen(discr, READ, BINARY);
do
    x:= nextgoodorder(id, n, f, roots);
    L:= [op(L), x];
    if x = NULL then
        break
    end if;
    i:= i + 2
end do;
fclose(discr);
L
end proc

```

```

> #
# This procedure test whether in the quadratic order with
# discriminant -D there is an element with norm equal to n,
# a probable prime. If such an element a+b*sqrt(-D) is found,
# the trace of it is given back, else the result is NULL.
#

```

```

testform:=proc(D,n) local a,b,r,a1;
if numtheory[jacobi](-D,n)<>1 then RETURN() fi;
b:=modsqrt(-D mod n,n);
if type(b+D,odd) then b:=b+n; fi;
r:=reduceform(n,b,(b^2+D)/4/n);
if D mod 4=0 then
    if r[1]<>1 or r[2]<>0 then RETURN() fi;
    2*r[4];
else
    if r[1]<>1 or r[2]<>1 then RETURN() fi;
    2*r[4]+r[5];
fi; end;

```

```

>

```

6. Factorization of the orders of curves

```

> ;

```

```

> ;

```

```

>

```

```
>
>
>
> ;
> ;
>
>
>
>
>
```

7. Set up the Hilbert polynomial

```
> #
# For a given negative discriminant -D this procedure
calculates
# a list of all pairs of numbers [a,b] representing a reduced
# ideal and gives back the list of these.
#
```

```
allredideals:=proc(D) local a,b,c,L; L:=[];
for a while 3*a^2<=D do
  for b from -a to a do
    if b^2+D mod 4*a <> 0 then next fi;
    c:=(b^2+D)/(4*a);
    if c<a then next fi;
    if igcd(a,b,c)>1 then next fi;
    if a=c and b<0 then next fi;
    if b=-a then next fi;
    L:=[op(L),[a,b]];
  od;
od; L; end;
```

allredideals := proc(D) (11.5.25)

```
local a, b, c, L;
L:= [];
for a while 3 * a^2 <= D do
  for b from -a to a do
    if mod(b^2 + D, 4 * a) <> 0 then
      next
    end if;
    c:= 1 / 4 * (b^2 + D) / a;
    if c < a then
      next
```

```

end if;
if  $l < \text{igcd}(a, b, c)$  then
  next
end if;
if  $a = c$  and  $b < 0$  then
  next
end if;
if  $b = -a$  then
  next
end if;
 $L := [\text{op}(L), [a, b]]$ 
end do
end do;
L
end proc

```

> #
This procedure calculate directly the parameters $g_2(L)$ and $g_3(L)$
of the elliptic curve over the complex numbers
corresponding
to a lattice $k+l*a_1$ with a complex number a_1 . The summing
is
done for all k, l with absolute value less than K . You may
change
Digits and increase K to obtain better approximation.
#

```

directparam:=proc(a1,K) local k,l,s4,s6; s6:=0;
for k while k<K do s4:=s4+2/k^4; s6:=s6+2/k^6; od;
for l while l<K do s4:=s4+2/(l*a1)^4; s6:=s6+2/(l*a1)^6; od;
for l while l<K do
  for k while k<K do
    s4:=s4+2/(k+l*a1)^4+2/(k-l*a1)^4;
    s6:=s6+2/(k+l*a1)^6+2/(k-l*a1)^6;
  od;
od; [60*s4,140*s6]; end;
directparam:=proc(a1,K)

```

(11.5.26)

```

local k, l, s4, s6;
s6:= 0;

```

```

for k while k < K do
    s4 := s4 + 2 / k4;
    s6 := s6 + 2 / k6
end do;
for l while l < K do
    s4 := s4 + 2 / (l4 * a4);
    s6 := s6 + 2 / (l6 * a6)
end do;
for l while l < K do
    for k while k < K do
        s4 := s4 + 2 / (k + l * a)4 + 2 / (k - l * a)4;
        s6 := s6 + 2 / (k + l * a)6 + 2 / (k - l * a)6
    end do
end do;
[60 * s4, 140 * s6]
end proc

> #
# This procedure calculate a crude approximation of the j-
# invariant
# corresponding to a reduced ideal with parameters a,b,-D.
# Simple summing is done for all k,l with absolute value less
# then K. You may change Digits and increase K to obtain
# better
# approximation.
#

directjinvariant := proc(a, b, D, K) local a1, p;
a1 := evalf((b + sqrt(-D)) / (2. * a));
p := directparam(a1, K);
26 * 33 * p[1]3 / (p[1]3 - 27 * p[2]2);
end;
directjinvariant := proc(a, b, D, K) (11.5.27)
    local a1, p;
    a1 := evalf((b + sqrt(-D)) / (2. * a));
    p := directparam(a1, K);
    1728 * p[1]3 / (p[1]3 - 27 * p[2]2)
end proc

```

```

> #
# This procedure calculates a crude approximation of the
# Hilbert polynomial for a given discriminant -D. To get the
# j-invariants, simple summing up for all k,l with absolute
# value less then K is used. You may change Digits and
# increase K
# to obtain better approximation.
#

directhilbert:=proc(D,K) local p,j,P,L; global X;
L:=allredideals(D);P:=1;
for p in L do
  j:=directjinvariant(p[1],p[2],D,K);
  P:=P*(X-j);
od; P; end;
directhilbert:= proc(D, K)
local p, j, P, L;
global X;
L:= allredideals(D);
P:= 1;
for pin Ldo
  j:= directjinvariant(p[1], p[2], D, K);
  P:= P* (X - j)
end do;
P
end proc

```

(11.5.28)

```

> #
# This procedure calculate the coefficients of the product of
# two pover series in order below n. The coefficients are
# given
# as lists starting with the constant term and ending with
# the
# term with order n-1.
#

cauchyprod:=proc(A,B,n) local c,i,j,C; C:=[];
for i from 0 to n-1 do c:=0;
  for j from 0 to i do c:=c+A[j+1]*B[i-j+1]; od;
  C:=[op(C),c];
od; C; end;
cauchyprod:= proc(A, B, n)
local c, i, j, C;

```

(11.5.29)

```

C:= [];
for i from 0 to n - 1 do
  c:= 0;
  for j from 0 to i do
    c:= c + A[j + 1] * B[i - j + 1]
  end do;
  C:= [op(C), c]
end do;
C
end proc

> #
# This procedure calculate the coefficients of the quotient
# of
# two power series in order below n. The coefficients are
# given
# as lists starting with the constant term and ending with
# the
# term with order n-1.
#

cauchyquo:=proc(A,B,n) local c,i,j,C; C:=[A[1]/B[1]];
for i to n-1 do c:=0;
  for j to i do c:=c+C[j]*B[i-j+2]; od;
  C:=[op(C), (A[i+1]-c)/B[1]];
od; C; end;
cauchyquo:=proc(A, B, n)
local c, i, j, C;
C:= [A[1]/B[1]];
for i to n - 1 do
  c:= 0;
  for j to i do
    c:= c + C[j]*B[i - j + 2]
  end do;
  C:= [op(C), (A[i + 1] - c) / B[1]]
end do;
C
end proc

```

(11.5.30)

```
> #
# This procedure calculate the coefficients of the q-expansion
# with order below n. The 1/q term will not be given back.
#
```

```
qexpansion:=proc(n) local c,i,j,C,A,B;
A:=[1,240*sigma[3](i)$i=1..n+1];
B:=[1,-504*sigma[5](i)$i=1..n+1];
C:=cauchyprod(A,A,n+2);
A:=cauchyprod(A,C,n+2);
B:=cauchyprod(B,B,n+2);
C:=[];
for i to n+1 do C:=[op(C),A[i+1]-B[i+1]]; od;
C:=cauchyquo(A,C,n+1);
map(i->1728*i,[C[2..n+1]]);
end;
```

qexpansion := **proc**(*n*) (11.5.31)

```
local c, i, j, C, A, B;
A := [1, $(240 * numtheory:-σ[3](i), i = 1..n + 1)];
B := [1, $(-504 * numtheory:-σ[5](i), i = 1..n + 1)];
C := cauchyprod(A, A, n + 2);
A := cauchyprod(A, C, n + 2);
B := cauchyprod(B, B, n + 2);
C := [];
for i to n + 1 do
    C := [op(C), A[i + 1] - B[i + 1]]
end do;
C := cauchyquo(A, C, n + 1);
map(proc(i)
    option operator, arrow;
    1728 * i
end proc, [C[2..n + 1]])
end proc
```

```
> #
# This procedure calculate a j-invariant corresponding to a
# reduced ideal given by a, b, and D. The prescribed
precision
# is Dig decimal digits. The coefficients of the q-expansion
# have to be given by the global list ccoeff. The length of
this
```

```

# vector is doubled, if necessary.
#

jinvariant:=proc(a,b,D,Dig) local olddig,q,r,qq,eps,i; global
ccoeff;
olddig:=Digits; Digits:=Dig;
q:=evalf(exp(2*Pi*I*(b+sqrt(-D))/2./a));r:=1/q;
qq:=1.; eps:=10.^(-Dig);
for i while abs(ccoeff[i]*qq)>abs(r)*eps do
  r:=r+ccoeff[i]*qq; qq:=qq*q;
  if i=nops(ccoeff) then ccoeff:=qexpansion(2*i); fi;
od; Digits:=olddig; r; end;
jinvariant:=proc(a, b, D, Dig)
  local olddig, q, r, qq, eps, i;
  global ccoeff;
  olddig:= Digits;
  Digits:= Dig;
  q:= evalf(exp(2 * pi * I * (b + sqrt(-D)) / (2. * a)));
  r:= 1 / q;
  qq:= 1.;
  eps:= 10.^(-Dig);
  for i while abs(r) * eps < abs(ccoeff[i] * qq) do
    r:= r + ccoeff[i] * qq;
    qq:= qq * q;
    if i = nops(ccoeff) then
      ccoeff:= qexpansion(2 * i)
    end if
  end do;
  Digits:= olddig;
  r
end proc
> #
# This procedure calculate the Hilbert polynomial
# corresponding
# to a discriminant -D. After all reduced ideals are
# calculated
# a preliminary calculation is done to obtain a bound for the
# l_1-norm of the coefficient vector of the result and log
[10]

```

(11.5.32)


```

# of this + Dig decimal digits will be the precision used.
The
# resulting polynomial of the global variable X and an
estimate
# of the error are given back. The coefficients of the q-
expansion
# have to be given by the global list ccoeff. The length of
this
# list is doubled, if necessary.
#

```

```

hilbert:=proc(D,Dig) local olddig,DD,ll,L,i,H,HH,eps,c;
global ccoeff,X;
olddig:=Digits; Digits:=Dig;
DD:=evalf(sqrt(D)); ll:=0; L:=allredideals(D);
for i in L do ll:=ll+log[10.](abs(exp(Pi*DD/i[1]))+2101.);
od;
Digits:=Dig+ceil(ll); H:=1.;
for i in L do H:=expand(H*(X-jinvariant(i[1],i[2],D,Digits)))
; od;
HH:=0; eps:=0.;
for i from 0 to nops(L) do
  c:=coeff(H,X,i); HH:=HH+round(c)*X^i;
  if abs(c-round(c))>eps then eps:=abs(c-round(c)); fi;
od; Digits:=olddig; HH,eps; end;

```

hilbert := **proc**(*D*, *Dig*) (11.5.33)

```

local olddig, DD, ll, L, i, H, HH, eps, c
global ccoeff, X;
olddig := Digits;
Digits := Dig;
DD := evalf(sqrt(D));
ll := 0;
L := allredideals(D);
for i in L do
  ll := ll + log[10.](abs(exp( $\pi$  * DD / i[1])) + 2101.)
end do;
Digits := Dig + ceil(ll);
H := 1.;
for i in L do
  H := expand(H * (X - jinvariant(i[1], i[2], D, Digits)))
end do;

```

```

HH:= 0;
eps:= 0.;
for ifrom 0 to nops(L) do
  c:= coeff(H, X, i);
  HH:= HH + round(c) * X^i;
  if eps < abs(c - round(c)) then
    eps:= abs(c - round(c))
  end if
end do;
Digits:= olddig;
HH, eps
end proc

```

8. Factoring the Hilbert polynomial

```

> #
# This procedure find the factorization of the polynomial P
# of
# the variable X modulo the prime number p in the form of
# product
# P_i^i, where each P_i is square free. The list of P_i's
# is given back.
#

squarefreefactor:=proc(P,p) local Pm,D,PP,L,LL,Q,R,i,j;
global X;
Pm:=P mod p;
if degree(Pm)<=1 then RETURN([Pm]); fi;
PP:=diff(Pm,X) mod p;
if PP=0 then
  for i from 0 while p*i<=degree(Pm) do PP:=PP+coeff(Pm,X,i*
p)*X^i; od;
  LL:=squarefreefactor(PP,p); L:=[];
  for i to nops(LL) do L:=[op(L),0$j=1..p-1,LL[i]]; od;
  RETURN(L);
fi;
D:=Gcd(Pm,PP) mod p; if degree(D)=0 then RETURN([Pm]); fi;
Q:=Quo(Pm,D,X) mod p; PP:=Gcd(Q,D) mod p; Divide(Q,PP,'R')
mod p;
[R,op(squarefreefactor(D,p))];
end;
squarefreefactor:= proc(P, p)

```

(11.5.34)

```

local Pm, D, PP, L, LL, Q, R, i, j;
global X;
Pm := mod(P, p);
if degree(Pm) <= 1 then
    RETURN([Pm])
end if;
PP := mod(diff(Pm, X), p);
if PP = 0 then
    for i from 0 while p*i <= degree(Pm) do
        PP := PP + coeff(Pm, X, p*i) * X^i
    end do;
    LL := squarefreefactor(PP, p);
    L := [];
    for i to nops(LL) do
        L := [op(L), $(0, j = 1..p-1), LL[i]]
    end do;
    RETURN(L)
end if;
D := mod(Gcd(Pm, PP), p);
if degree(D) = 0 then
    RETURN([Pm])
end if;
Q := mod(Quo(Pm, D, X), p);
PP := mod(Gcd(Q, D), p);
mod(Divide(Q, PP, 'R'), p);
[R, op(squarefreefactor(D, p))]
end proc

```

```

> #
# This procedure try to find a root of a polynomial P of X
# modulo p. The polynomial is known to be product of degree
# 1 irreducible polynomials. The parameter K is the limit
# for the random tries. The root found or FAIL is given back.
#

```

```

modpolyroot := proc(P, p, K) local L, PP, PPP, i, r, T, D;

```

```

PP:=1; r:=rand(p); L:=squarefreefactor(P,p);
for i to nops(L) do
  PPP:=L[i];
  if degree(PPP)>0 then
    if degree(PP)=0 or degree(PP)>degree(PPP) then PP:=PPP;
  fi;
  fi;
od;
if degree(PP)=0 then RETURN(false) fi;
while degree(PP)>1 do
  for i to K do
    T:=0; T:=X+r(); T:=Powmod(T,(p-1)/2,PP,X) mod p; D:=Gcd
(P,T+1) mod p;
    if degree(D)>0 and degree(D)<degree(PP) then
      if degree(D)<=degree(PP)/2 then PP:=D; else PP:=Quo(PP,
D,X) mod p; fi;
      break;
    fi;
  od;
  if i>K then RETURN(FAIL) fi;
od; -coeff(PP,X,0)/coeff(PP,X,1) mod p; end;
modpolyroot:=proc(P, p, K)

```

(11.5.35)

```

local L, PP, PPP, i, r, T, D;
PP:= 1;
r:= rand(p);
L:= squarefreefactor(P, p);
for i to nops(L) do
  PPP:= L[i];
  if 0 < degree(PPP) then
    if degree(PP) = 0 or degree(PPP) < degree(PP) then
      PP:= PPP
    end if
  end if
end do;
if degree(PP) = 0 then
  RETURN(false)
end if;
while 1 < degree(PP) do
  for i to K do

```

```

T:= 0;
T:= X+ r();
T:= mod(Powmod(T, 1 / 2 * p - 1 / 2, PP, X), p);
D := mod(Gcd(P, T+ 1), p);
if 0 < degree(D) and degree(D) < degree(PP) then
    if degree(D) <= 1 / 2 * degree(PP) then
        PP:= D
    else
        PP:= mod(Quo(PP, D, X), p)
    end if;
    break
end if
end do;
if K < i then
    RETURN(FAIL)
end if
end do;
mod(-coeff(PP, X, 0) / coeff(PP, X, 1), p)
end proc

```

9. Set up the curves and the points

```

> #
# This procedure calculate a random point on a given elliptic
# curve
# modulo n with parameters a,b.
#
ellrandpoint:=proc(a,b,n) local i,j,x,y,r;
r:=rand(n); j:=-1;
for i to 100 while j<>1 do
    x:=r(); j:=numtheory[jacobi](x^3+a*x+b,n);
od;
if i=100 then ERROR(`probably not a prime`,n) fi;
y:=modsqrt(x^3+a*x+b mod n,n); j:=r();
if j>n/2 then [x,y,1]; else [x,-y mod n,1]; fi;
end;
ellrandpoint:= proc(a, b, n)

```

(11.5.36)

```

local i, j, x, y, r;
r := rand(n);
j := -1;
for i to 100 while j <> 1 do
    x := r();
    j := numtheory[jacobi](x^3 + a*x + b, n)
end do;
if i = 100 then
    ERROR(probably not a prime, n)
end if;
y := modsqrt(mod(x^3 + a*x + b, n), n);
j := r();
if 1 / 2 * n < j then
    [x, y, 1]
else
    [x, mod(-y, n), 1]
end if
end proc

```

```
>
```

```
>
```

10. Putting all together

```
> interface(echo=3);
2 (11.5.37)
```

```
> ;
```

```
> #
```

```
# This simple Maple routine calculates the function
#  $L(x, a, b) = \exp(b \cdot (\ln(x))^a \cdot (\ln(\ln(x)))^{(1-a)})$ 
#
```

```
L:=proc(x, a, b) evalf(exp(b*(ln(x))^a*(ln(ln(x)))^(1-a))) end;
L:= proc(x, a, b) (11.5.38)
```

```
evalf(exp(b*ln(x)^a*ln(ln(x))^(1 - a)))
```

```
end proc
```

```
> #
```

Calculation of running time estimates

#

> **D2:=n->m(ln(n))*d(n)/ln(d(n))*rep(n);**

$$D2 := n \rightarrow \frac{m(\ln(n)) d(n) \text{rep}(n)}{\ln(d(n))} \quad (11.5.39)$$

> **D3:=n->m(ln(n))*ln(ln(n))*d(n)/ln(d(n))*rep(n);**

$$D3 := n \rightarrow \frac{m(\ln(n)) \ln(\ln(n)) d(n) \text{rep}(n)}{\ln(d(n))} \quad (11.5.40)$$

> **aa:=n->e(n)*b(n)*ln(n)*rep(n);**

$$aa := n \rightarrow e(n) b(n) \ln(n) \text{rep}(n) \quad (11.5.41)$$

> **bb:=n->ln(e(n)*ln(n))*b(n)/e(n)/ln(n)*m(e(n)*ln(n))*rep(n);**

$$bb := n \rightarrow \frac{\ln(e(n) \ln(n)) b(n) m(e(n) \ln(n)) \text{rep}(n)}{e(n) \ln(n)} \quad (11.5.42)$$

> **cc:=n->0;**

$$cc := n \rightarrow 0 \quad (11.5.43)$$

> **dd:=n->e(n)*b(n)^(1/2)*m(ln(n))*rep(n);**

$$dd := n \rightarrow e(n) \sqrt{b(n)} m(\ln(n)) \text{rep}(n) \quad (11.5.44)$$

> **ee:=n->e(n)*m(ln(n))*L(b(n),sqrt(2))*rep(n);**

$$ee := n \rightarrow e(n) m(\ln(n)) L(b(n), \sqrt{2}) \text{rep}(n) \quad (11.5.45)$$

> **D6:=n->e(n)*ln(n)*m(ln(n))*rep(n);**

$$D6 := n \rightarrow e(n) \ln(n) m(\ln(n)) \text{rep}(n) \quad (11.5.46)$$

> **A:=n->m(h(DD)*ln(n))*ln(n)*rep(n);**

$$A := n \rightarrow m(h(DD) \ln(n)) \ln(n) \text{rep}(n) \quad (11.5.47)$$

> **B:=n->m(g(DD)*ln(n))*ln(n)*rep(n);**

$$B := n \rightarrow m(g(DD) \ln(n)) \ln(n) \text{rep}(n) \quad (11.5.48)$$

> **m:=x->x*ln(x)*ln(ln(x));**

$$m := x \rightarrow x \ln(x) \ln(\ln(x)) \quad (11.5.49)$$

> **L:=(x,beta)->exp(beta*((ln(x)*ln(ln(x)))^(1/2)));**

$$L := (x, \beta) \rightarrow e^{\beta \sqrt{\ln(x) \ln(\ln(x))}} \quad (11.5.50)$$

> **rep:=n->ln(n)/ln(b(n));**

$$\text{rep} := n \rightarrow \frac{\ln(n)}{\ln(b(n))} \quad (11.5.51)$$

> **h:=x->x^(1/2);**

$$h := x \rightarrow \sqrt{x} \quad (11.5.52)$$

> **g:=x->x^(1/2)/ln(x);**

$$(11.5.53)$$

$$g := x \rightarrow \frac{\sqrt{x}}{\ln(x)} \quad (11.5.53)$$

> DD:=d(n);

$$DD := \ln(n)^2 \quad (11.5.54)$$

> d:=x->(ln(x))^2; e:=x->(d(x))^(1/2); b:=x->ln(x);

$$d := x \rightarrow \ln(x)^2$$

$$e := x \rightarrow \sqrt{d(x)}$$

$$b := x \rightarrow \ln(x) \quad (11.5.55)$$

> expand(D2(n));
 expand(D3(n));
 expand(aa(n));
 expand(D6(n));
 expand(A(n));
 expand(B(n));

$$\frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)}$$

$$\frac{\ln(n)^4 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3}{\ln(\ln(n))}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\ln(\ln(n)))}{\ln(\ln(n))}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\sqrt{\ln(n)^2} \ln(n)) \ln(\ln(\sqrt{\ln(n)^2} \ln(n)))}{\ln(\ln(n))}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln(\ln(n)^2) \ln(\ln(n))} \quad (11.5.56)$$

> d:=x->(ln(x))^2; e:=x->(d(x))^(1/2); b:=x->ln(x);

$$d := x \rightarrow \ln(x)^2$$

$$e := x \rightarrow \sqrt{d(x)}$$

$$b := x \rightarrow \ln(x) \quad (11.5.57)$$

> expand(D2(n));
 expand(D3(n));
 expand(aa(n));
 expand(D6(n));
 expand(A(n));
 expand(B(n));

$$\begin{aligned}
& \frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)} \\
& \frac{\ln(n)^4 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^3}{\ln(\ln(n))} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\sqrt{\ln(n)^2} \ln(n)) \ln(\ln(\sqrt{\ln(n)^2} \ln(n)))}{\ln(\ln(n))} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln(\ln(n)^2) \ln(\ln(n))}
\end{aligned} \tag{11.5.58}$$

> **d:=x->(ln(x))^2; e:=x->(d(x))^(1/2); b:=x->(ln(x))^2;**
d:= x→ln(x)²

e:= x→√d(x)

b:= x→ln(x)²

(11.5.59)

> **expand(D2(n));**
expand(D3(n));
expand(aa(n));
expand(D6(n));
expand(A(n));
expand(B(n));

$$\begin{aligned}
& \frac{\ln(n)^4 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)^2} \\
& \frac{\ln(n)^4 \ln(\ln(n))^2 \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)^2} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^4}{\ln(\ln(n)^2)} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\sqrt{\ln(n)^2} \ln(n)) \ln(\ln(\sqrt{\ln(n)^2} \ln(n)))}{\ln(\ln(n)^2)}
\end{aligned}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln(\ln(n)^2)^2} \quad (11.5.60)$$

> **b:=x->(ln(x))^2; e:=x->(d(x))^(1/2);**
b:=x->ln(x)^(ln(ln(x))*ln(ln(ln(x))))^(-2));
b:=x->ln(x)^2

$$e:=x \rightarrow \sqrt{d(x)}$$

$$b:=x \rightarrow \ln(x) \frac{\ln(\ln(x))}{\ln(\ln(\ln(x)))^2} \quad (11.5.61)$$

> **expand(D2(n));**
expand(D3(n));
expand(aa(n));
expand(D6(n));
expand(A(n));
expand(B(n));

$$\frac{\ln(n)^4 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln(\ln(n)^2) \ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)}$$

$$\frac{\ln(n)^4 \ln(\ln(n))^2 \ln(\ln(\ln(n)))}{\ln(\ln(n)^2) \ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2} \ln(n)^2}{\ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln(\ln(n)^2) \ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)} \quad (11.5.62)$$

> **d:=x->(ln(x))^2/(ln(ln(x)))^2; e:=x->(d(x))^(1/2); b:=x->ln(x);**

$$d := x \rightarrow \frac{\ln(x)^2}{\ln(\ln(x))^2}$$

$$e := x \rightarrow \sqrt{d(x)}$$

$$b := x \rightarrow \ln(x)$$

(11.5.63)

> **expand(D2(n));**
expand(D3(n));
expand(aa(n));
expand(D6(n));
expand(A(n));
expand(B(n));

$$\frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln(\ln(n))^2 \ln\left(\frac{\ln(n)^2}{\ln(\ln(n))^2}\right)}$$

$$\frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln(\ln(n)) \ln\left(\frac{\ln(n)^2}{\ln(\ln(n))^2}\right)}$$

$$\frac{\sqrt{\frac{\ln(n)^2}{\ln(\ln(n))^2}} \ln(n)^3}{\ln(\ln(n))}$$

$$\sqrt{\frac{\ln(n)^2}{\ln(\ln(n))^2}} \ln(n)^3 \ln(\ln(\ln(n)))$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\sqrt{\ln(n)^2} \ln(n)) \ln(\ln(\sqrt{\ln(n)^2} \ln(n)))}{\ln(\ln(n))}$$

$$\frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln(\ln(n)^2) \ln(\ln(n))}$$

(11.5.64)

> **d:=x->(ln(x))^2/(ln(ln(x)))^2; e:=x->(d(x))^(1/2);**
b:=x->(ln(x))^3/ln(ln(n))^3;

$$d := x \rightarrow \frac{\ln(x)^2}{\ln(\ln(x))^2}$$

$$e := x \rightarrow \sqrt{d(x)}$$

$$b := x \rightarrow \frac{\ln(x)^3}{\ln(\ln(n))^3}$$

(11.5.65)

> **expand(D2(n));**
expand(D3(n));

expand(aa(n));
expand(D6(n));
expand(A(n));
expand(B(n));

$$\begin{aligned}
 & \frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln(\ln(n)) \ln\left(\frac{\ln(n)^2}{\ln(\ln(n))^2}\right) \ln\left(\frac{\ln(n)^3}{\ln(\ln(n))^3}\right)} \\
 & \frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln\left(\frac{\ln(n)^2}{\ln(\ln(n))^2}\right) \ln\left(\frac{\ln(n)^3}{\ln(\ln(n))^3}\right)} \\
 & \frac{\sqrt{\frac{\ln(n)^2}{\ln(\ln(n))^2}} \ln(n)^5}{\ln(\ln(n))^3 \ln\left(\frac{\ln(n)^3}{\ln(\ln(n))^3}\right)} \\
 & \frac{\sqrt{\frac{\ln(n)^2}{\ln(\ln(n))^2}} \ln(n)^3 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln\left(\frac{\ln(n)^3}{\ln(\ln(n))^3}\right)} \\
 & \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln(\sqrt{\ln(n)^2} \ln(n)) \ln(\ln(\sqrt{\ln(n)^2} \ln(n)))}{\ln\left(\frac{\ln(n)^3}{\ln(\ln(n))^3}\right)} \\
 & \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln(\ln(n)^2) \ln\left(\frac{\ln(n)^3}{\ln(\ln(n))^3}\right)}
 \end{aligned} \tag{11.5.66}$$

> d:=x->(ln(x))^2/(ln(ln(x)))^2; e:=x->(d(x))^(1/2); b:=x->(ln(x))^2;

$$d := x \rightarrow \frac{\ln(x)^2}{\ln(\ln(x))^2}$$

$$e := x \rightarrow \sqrt{d(x)}$$

$$b := x \rightarrow \ln(x)^2$$

(11.5.67)

> expand(D2(n));
expand(D3(n));
expand(aa(n));
expand(D6(n));
expand(A(n));

expand(B(n)) ;

$$\begin{aligned}
 & \frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln(\ln(n)) \ln\left(\frac{\ln(n)^2}{\ln(\ln(n))^2}\right) \ln(\ln(n)^2)} \\
 & \frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln\left(\frac{\ln(n)^2}{\ln(\ln(n))^2}\right) \ln(\ln(n)^2)} \\
 & \frac{\sqrt{\frac{\ln(n)^2}{\ln(\ln(n))^2}} \ln(n)^4}{\ln(\ln(n)^2)} \\
 & \frac{\sqrt{\frac{\ln(n)^2}{\ln(\ln(n))^2}} \ln(n)^3 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln(\ln(n)^2)} \\
 & \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\sqrt{\ln(n)^2} \ln(n)\right) \ln\left(\ln\left(\sqrt{\ln(n)^2} \ln(n)\right)\right)}{\ln(\ln(n)^2)} \\
 & \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln(\ln(n)^2)^2}
 \end{aligned} \tag{11.5.68}$$

**> d:=x->(ln(x))^2/(ln(ln(x)))^2; e:=x->(d(x))^(1/2);
b:=x->ln(x)^(ln(ln(x))*ln(ln(ln(x))))^(-2));**

$$\begin{aligned}
 d &:= x \rightarrow \frac{\ln(x)^2}{\ln(\ln(x))^2} \\
 e &:= x \rightarrow \sqrt{d(x)} \\
 b &:= x \rightarrow \ln(x) \frac{\ln(\ln(x))}{\ln(\ln(\ln(x)))^2}
 \end{aligned} \tag{11.5.69}$$

**> expand(D2(n));
expand(D3(n));
expand(aa(n));
expand(D6(n));
expand(A(n));
expand(B(n));**

$$\frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln(\ln(n)) \ln\left(\frac{\ln(n)^2}{\ln(\ln(n))^2}\right) \ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)}$$

$$\begin{aligned}
& \frac{\ln(n)^4 \ln(\ln(\ln(n)))}{\ln\left(\frac{\ln(n)^2}{\ln(\ln(n))^2}\right) \ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)} \\
& \frac{\sqrt{\frac{\ln(n)^2}{\ln(\ln(n))^2}} \ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2} \ln(n)^2}{\ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)} \\
& \frac{\sqrt{\frac{\ln(n)^2}{\ln(\ln(n))^2}} \ln(n)^3 \ln(\ln(n)) \ln(\ln(\ln(n)))}{\ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\sqrt{\ln(n)^2} \ln(n)\right) \ln\left(\ln\left(\sqrt{\ln(n)^2} \ln(n)\right)\right)}{\ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)} \\
& \frac{\sqrt{\ln(n)^2} \ln(n)^3 \ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right) \ln\left(\ln\left(\frac{\sqrt{\ln(n)^2} \ln(n)}{\ln(\ln(n)^2)}\right)\right)}{\ln(\ln(n)^2) \ln\left(\ln(n) \frac{\ln(\ln(n))}{\ln(\ln(\ln(n)))^2}\right)}
\end{aligned} \tag{11.5.70}$$



```

> #
# The following procedure has as input a list
# [n,a,b,x,y,s,n1,a1,b1,x1,y1,s1,...,nn], and checks whether
# this list represents a "partial proof" of the primality
# of n, i.e. a proof with elliptic curves that if nn is a
# prime then n is a prime, too. The steps are to check
whether
# n_i is a prime if n_{i+1} is a prime. This consist of
cheking
# that a_i, b_i are the parameters of an elliptic curve
# modulo n_i, the point P=(x_i,y_i,1) is on this elliptic
# curve, that Q=(m_i/n_{i+1})*P<>0 and n_{i+1}*Q=0, the zero
# of the elliptic curve. Here m_i, the cardinality of the
# elliptic curve is n_{i+1}+s_i.
#

```

```

partialproof:=proc(PP) local i,n,a,b,P,Q,x,y,m,np;
if not type(LL,list) then RETURN(false) fi;
if nops(PP) mod 6<>1 then RETURN(false) fi;
for i to nops(LL) do
  if not type(LL[i],nonnegint) then RETURN(false) fi;
od;
for i by 6 while i<nops(PP) do
  n:=PP[i]; if gcd(n,6)>1 then RETURN(false); fi;
  a:=PP[i+1]; b:=PP[i+2];
  if gcd(n,4*a^3+27*b^2)>1 then RETURN(false); fi;
  x:=PP[i+3]; y:=PP[i+4];
  if y^2-x^3-a*x-b mod n<>0 then RETURN(false) fi;
  P:=[x,y,1]; m:=n+1+PP[i+5]; np:=PP[i+6];
  P:=ellmul(P,m/np,a,b,n);
  if P[3]<>1 then RETURN(false); fi;
  P:=ellmul(P,np,a,b,n);
  if P[3]<>0 then RETURN(false); fi;
od; true; end;

```

partialproof:= proc(PP) (11.5.71)

```

local i, n, a, b, P, Q, x, y, m, np;
if not type(LL, list) then
  RETURN(false)
end if;
if mod(nops(PP), 6) <> 1 then
  RETURN(false)
end if;
for i to nops(LL) do

```

```

if not type(LL[i], nonnegint) then
    RETURN(false)
end if
end do;
for i by 6 while i < nops(PP) do
    n := PP[i];
    if 1 < gcd(n, 6) then
        RETURN(false)
    end if;
    a := PP[i + 1];
    b := PP[i + 2];
    if 1 < gcd(n, 4 * a^3 + 27 * b^2) then
        RETURN(false)
    end if;
    x := PP[i + 3];
    y := PP[i + 4];
    if mod(y^2 - x^3 - a * x - b, n) <> 0 then
        RETURN(false)
    end if;
    P := [x, y, 1];
    m := n + 1 + PP[i + 5];
    np := PP[i + 6];
    P := ellmul(P, m / np, a, b, n);
    if P[3] <> 1 then
        RETURN(false)
    end if;
    P := ellmul(P, np, a, b, n);
    if P[3] <> 0 then
        RETURN(false)
    end if
end do;
true

```


end proc

```
> #  
# The following procedure has as input a list  
# [n,a,b,x,y,s,n1,a1,b1,x1,y1,s1,...,nn], and checks whether  
# this list represents a proof of the primality of n with  
# elliptic curves. First the primality of nn checked with  
# trial division. Then the procedure partialproof is called  
# to complete the proof.  
#
```

```
proof:=proc(PP) local i,nn;  
if not type(LL,list) then RETURN(false) fi;  
if nops(PP) mod 6<>1 then RETURN(false) fi;  
if not type(LL[nops(LL)],nonnegint) then RETURN(false) fi;  
nn:=PP[nops(PP)];  
if type(nn,even) then RETURN(false) fi;  
for i from 3 by 2 while i*i<=nn do  
  if nn mod i=0 then RETURN(false) fi;  
od; partialproof(PP); end;
```

proof:= proc(PP)

(11.5.72)

```
local i, nn;  
if not type(LL, list) then  
  RETURN(false)  
end if;  
if mod(nops(PP), 6) <> 1 then  
  RETURN(false)  
end if;  
if not type(LL[nops(LL)], nonnegint) then  
  RETURN(false)  
end if;  
nn := PP[nops(PP)];  
if type(nn, even) then  
  RETURN(false)  
end if;  
for i from 3 by 2 while i*i <= nn do  
  if mod(nn, i) = 0 then  
    RETURN(false)  
  end if
```

```
end do;  
  partialproof(PP)  
end proc
```

- ▶ **12. Polinomfaktorizálás**
- ▶ **13. Az AKS-teszt**
- ▶ **14. A szita módszerek alapjai**
- ▶ **15. Számtest szita**
- ▶ **16. Vegyes problémák**