

Számítógépes számelmélet

Járai Antal

Ezek a programok csak szemléltetésre szolgálnak

- ▶ 1. A prímek eloszlása, szitálás
- ▶ 2. Egyszerű faktorizálási módszerek
- ▶ 3. Egyszerű prímtesztelési módszerek
- ▶ 4. Lucas-sorozatok
- ▶ 5. Alkalmazások
- ▶ 6. Számok és polinomok
- ▶ 7. Gyors Fourier-transzformáció
- ▶ 8. Elliptikus függvények
- ▶ 9. Számolás elliptikus görbéken
- ▼ 10. Faktorizálás elliptikus görbékkel

```
> restart;
> #
# This routine randomly choose an elliptic "curve" modulo n,
# where gcd(n,6)=1. The coordinates x,y are chosen
# randomly, the parameter a too, and b is calculated.
# The list [x,y,a,b] is given back or a divisor d of n.
#

ellrand:=proc(n) local x,y,a,b,r,d,f;
r:=rand(n); d:=n;
while d=n do
  x:=r(n); y:=r(n); a:=r(n); b:=y^2-x^3-a*x mod n;
  d:=4*a^3+27*b^2 mod n; gcd(d,n);
end do;
end proc;
```

```

od;
if %<n and %>1 then return % fi;
[x,y,a,b]; end;
ellrand:=proc(n)
local x, y, a, b, r, d, f;
r:=rand(n);
d:=n;
while d = ndo
x:=r(n);
y:=r(n);
a:=r(n);
b:=mod(y^2 - x^3 - a*x, n);
d:=mod(4*a^3 + 27*b^2, n);
gcd(d, n)
end do;
if% < n and 1 < % then
return %
end if;
[x, y, a, b]
end proc

```

(10.1)

```

> #
# Addition on an elliptic "curve" modulo n, where gcd(n,6)=1.
# P and Q are the points to add and a,b are the parameters.
# The return value is the sum of the points or a divisor d of n.
#

```

```

elladd:=proc(P,Q,a,b,n) local l,d;
if P[3]=0 then return Q fi;
if Q[3]=0 then return P fi;
if P=Q then return elldou(P,a,b,n) fi;
if P[1]=Q[1] then return [0,1,0] fi;
d:=igcdex(P[1]-Q[1],n,'l');
if 1<d and d<n then return d fi;
l:=(P[2]-Q[2])*l mod n;
l^2-P[1]-Q[1] mod n;
[%,l*(P[1]-%)-P[2] mod n,1];
end;
elladd:=proc(P, Q, a, b, n)
local l, d;
if P[3] = 0 then
return Q
end if;

```

(10.2)

```

if Q[3] = 0 then
    return P
end if;
if P = Q then
    return elldou(P, a, b, n)
end if;
if P[1] = Q[1] then
    return [0, 1, 0]
end if;
d := igcdex(P[1] - Q[1], n, 'l');
if 1 < d and d < n then
    return d
end if;
l := mod((P[2] - Q[2])*l, n);
mod(l^2 - P[1] - Q[1], n);
[%, mod(l*(P[1] - %) - P[2], n), 1]
end proc

```

```

> #
# Doubling on an elliptic "curve" modulo n, where gcd(n,6)=1.
# P is the point to double and a, b are the parameters.
# The return value is the double of the point P or
# a divisor d of n.
#

```

```

elldou:=proc(P,a,b,n) local l,d;
if P[3]=0 then return P fi;
if P[2]=0 then return [0,1,0] fi;
d:=igcdex(2*P[2],n,'l');
if 1<d and d<n then return d fi;
l:=(3*P[1]^2+a)*l mod n;
l^2-2*P[1] mod n;
[%,l*(P[1]-%)-P[2] mod n,1];
end;

```

```

elldou:=proc(P, a, b, n)

```

(10.3)

```

local l, d;
if P[3] = 0 then
    return P
end if;
if P[2] = 0 then
    return [0, 1, 0]
end if;

```

```

d:= igcdex(2*P[2], n, T);
if 1 < d and d < n then
    return d
end if;
l:= mod((3*P[1]^2 + a)*l, n);
mod(l^2 - 2*P[1], n);
[%, mod(l*(P[1] - %) - P[2], n), 1]
end proc

```

```

> #
# This program compute k*P, k>=0 or a divisor of n
# on an elliptic "curve" modulo n, where gcd(n,6)=1.
# It use the left-to-right binary method.
#

```

```

ellmul:=proc(P,k,a,b,n) local L,i,Q;
if k=0 then return [0,1,0] fi;
if P[3]=0 then return P fi;
L:=convert(k,base,2);
Q:=P;
for i from nops(L)-1 to 1 by -1 do
    Q:=elldou(Q,a,b,n);
    if type(Q,integer) then return Q fi;
    if L[i]=1 then
        Q:=elladd(P,Q,a,b,n);
        if type(Q,integer) then return Q fi;
    fi;
od; Q; end;

```

```

ellmul:=proc(P, k, a, b, n)

```

(10.4)

```

local L, i, Q;
if k = 0 then
    return [0, 1, 0]
end if;
if P[3] = 0 then
    return P
end if;
L:= convert(k, base, 2);
Q:= P;
for i from nops(L) - 1 by -1 to 1 do
    Q:= elldou(Q, a, b, n);
    if type(Q, integer) then
        return Q
    end if;
end if;

```

```

    if L[i] = 1 then
        Q := elladd(P, Q, a, b, n);
        if type(Q, integer) then
            return Q
        end if
    end if
end do;
Q
end proc
> #
# This program compute k*P on an elliptic curve
# with parameters a,b mod n, where k=k(s,B) is the product of
# all prime-powers p^ep for which p<=s and p^ep<=B.
# B is roughly the bound for the prime factor we
# try to found. Usually s is some hundred and B is some million.
# If a divisor d of n found then d is given back.
# We suppose that gcd(n,6)=1.
#

elltry:=proc(P,s,B,a,b,n) local lB,p,ep,Q;
Q:=P; lB:=evalf(ln(B));
for p from 2 to s do
    if isprime(p) then
        ep:=floor(lB/evalf(ln(p)));
        Q:=ellmul(Q,p^ep,a,b,n);
        if type(Q,integer) then return Q fi;
    fi;
od; Q; end;
elltry:=proc(P,s,B,a,b,n)
local lB,p,ep,Q;
Q:=P;
lB:=evalf(ln(B));
for p from 2 to s do
    if isprime(p) then
        ep:=floor(lB/evalf(ln(p)));
        Q:=ellmul(Q,p^ep,a,b,n);
        if type(Q,integer) then
            return Q
        end if
    end if
end do;
Q

```

(10.5)

end proc

```
> #  
# This program try to split an integer n with the  
# elliptic curve method. The parameters are chosen to be  
# optimal to find prime factors below  
# P with probability 1-p. Better to remove first  
# small prime divisors. A factor found is given back.  
#
```

```
ellsplit:=proc(n,P,p) local d,B,s,K,i,a,b,Q;  
if isprime(n) then RETURN([n]) fi;  
B:=ceil(P+2.*sqrt(P)+1);  
s:=evalf(exp(sqrt(ln(P)*ln(ln(P))/2)));  
s:=ceil(max(3,s));  
K:=ceil(max(3,-s*ln(p)));  
for i to K do;  
  d:=ellrand(n);  
  if type(d,integer) then break fi;  
  Q:=[d[1],d[2],1]; a:=d[3]; b:=d[4];  
  d:=elltry(Q,s,B,a,b,n);  
  if type(d,integer) then break fi;  
od;  
if type(d,integer) then d else 1 fi; end;
```

```
ellsplit:=proc(n,P,p)
```

(10.6)

```
local d, B, s, K, i, a, b, Q;
```

```
if isprime(n) then
```

```
  RETURN([n])
```

```
end if;
```

```
B:=ceil(P+2.*sqrt(P)+1);
```

```
s:=evalf(exp(sqrt(1/2*ln(P)*ln(ln(P)))));
```

```
s:=ceil(max(3,s));
```

```
K:=ceil(max(3,-s*ln(p)));
```

```
for i to K do
```

```
  d:=ellrand(n);
```

```
  if type(d,integer) then
```

```
    break
```

```
  end if;
```

```
  Q:=[d[1],d[2],1];
```

```
  a:=d[3];
```

```
  b:=d[4];
```

```
  d:=elltry(Q,s,B,a,b,n);
```

```
  if type(d,integer) then
```

```

        break
    end if
end do;
if type(d, integer) then
    d
else
    1
end if
end proc
> ellsplit(1001,20,0.001);

```

13

(10.7)

▼ 10.1. Kötegelt végrehajtás.

```

> #
# Batch modular inverse modulo n for all L[i] or a divisor of
# n.
#

batchmodinv:=proc(L,n) local i,l,d,LL,LLL;
if nops(L)=1 then
    d:=igcdex(L[1],n,'1');
    if 1<d then return d fi;
    [];
else
    LL:=[]; i:=1;
    while i<=nops(L) do
        if i=nops(L) then
            LL:=[op(LL),L[i] mod n]
        else
            LL:=[op(LL),L[i]*L[i+1] mod n]
        fi; i:=i+2;
    od;
    LL:=batchmodinv(LL,n);
    if type(LL,integer) then return LL fi;
    while i<=nops(L) do
        if i=nops(L) then
            LL:=[op(LL),L[i] mod n]
        else
            LL:=[op(LL),L[i]*L[i+1] mod n]
        fi; i:=i+2;
    od;
    LLL:=[]; i:=1;
    while i<=nops(L) do
        if i=nops(L) then

```

```

        LLL:=[op(LL),LL[(i-1)/2]]
    else
        LLL:=[op(LL),L[i+1]*LL[(i-1)/2] mod n,L[i]*LL[(i-1)/2]
mod n]
        fi; i:=i+2;
    od;
    LLL;
fi; end;
batchmodinv:=proc(L, n)
local i, l, d, LL, LLL;
if nops(L) = 1 then
    d:=igcdex(L[1], n, 'T');
    if 1 < d then
        return d
    end if;
    [l]
else
    LL:=[];
    i:=1;
    while i <= nops(L) do
        if i = nops(L) then
            LL:= [op(LL), mod(L[i], n)]
        else
            LL:= [op(LL), mod(L[i]*L[i+1], n)]
        end if;
        i:= i + 2
    end do;
    LL:= modinvbatch(LL, n);
    if type(LL, integer) then
        return LL
    end if;
    while i <= nops(L) do
        if i = nops(L) then
            LL:= [op(LL), mod(L[i], n)]
        else
            LL:= [op(LL), mod(L[i]*L[i+1], n)]
        end if;
        i:= i + 2
    end do;

```

(10.1.1)


```

LLL:= [];
i:= 1;
while i <= nops(L) do
  if i = nops(L) then
    LLL:= [op(LLL), LL[1/2*i - 1/2]]
  else
    LLL:= [op(LLL), mod(L[i+1]*LL[1/2*i - 1/2],
n), mod(L[i]*LL[1/2*i - 1/2], n)]
  end if;
  i:= i + 2
end do;
LLL
end if
end proc

```

▼ 10.2. Szorzás egészekkel.

```

> #
# Doubling on an elliptic "curve" modulo n, where gcd(n,6)=1.
# x is the first coordinate of the point to double and a, b
# are the parameters. The return value is either a list, the
# first
# coordinate of the double of the point or a divisor d of n.
#

ellxdou:=proc(x,a,b,n) local l,d;
d:=igcdex(4*(x^3+a*x+b),n,'l');
if 1<d then d else [l*((x^2-a)^2-8*b*x) mod n] fi;
end;
ellxdou:=proc(x, a, b, n)
local l, d;
d:=igcdex(4*x^3 + 4*a*x + 4*b, n, 'l');
if 1 < d then
d
else
[mod(l*((x^2 - a)^2 - 8*b*x), n)]
end if
end proc
(10.2.1)

> #
# Calculation of the triple of first coordinates
# [x,x_2i,x_2i+1] if t=0 or [x,x_2i+1,x_2i+2] if t=1

```

```

# from the triple of first coordinates [x,x_i,x_{i+1}]
# on an elliptic "curve" modulo n, where gcd(n,6)=1, and a, b
# are the parameters. The return value is this list
# or a divisor d of n.
#

```

```

ellnextx:=proc(L,t,a,b,n) local l,d,x,xi,xip,xii,xiip;
x:=L[1]; xi:=L[2]; xip:=L[3];
if t=0 then
  d:=igcdex(4*(xi^3+a*xi+b),n,'1');
  if 1<d then return d fi;
  xii:=1*((xi^2-a)^2-8*b*xi) mod n;
  d:=igcdex(x*(xi-xip)^2,n,'1');
  if 1<d then return d fi;
  xiip:=1*((a-xi*xip)^2-4*b*(xi+xip)) mod n;
else
  d:=igcdex(x*(xi-xip)^2,n,'1');
  if 1<d then return d fi;
  xii:=1*((a-xi*xip)^2-4*b*(xi+xip)) mod n;
  d:=igcdex(4*(xip^3+a*xip+b),n,'1');
  if 1<d then return d fi;
  xiip:=1*((xip^2-a)^2-8*b*xip) mod n;
fi; [x,xii,xiip] end;

```

ellnextx := **proc**(*L*, *t*, *a*, *b*, *n*)

(10.2.2)

local *l*, *d*, *x*, ξ , *xip*, *xii*, *xiip*;

x := *L*[1];

ξ := *L*[2];

xip := *L*[3];

if *t* = 0 **then**

d := *igcdex*($4 * \xi^3 + 4 * a * \xi + 4 * b$, *n*, 'T');

if 1 < *d* **then**

return *d*

end if;

xii := *mod*($1 * ((\xi^2 - a)^2 - 8 * b * \xi)$, *n*);

d := *igcdex*($x * (\xi - xip)^2$, *n*, 'T');

if 1 < *d* **then**

return *d*

end if;

xiip := *mod*($1 * ((a - \xi * xip)^2 - 4 * b * (\xi + xip))$, *n*)

else

d := *igcdex*($x * (\xi - xip)^2$, *n*, 'T');

if 1 < *d* **then**

return *d*

```

end if;
xii:= mod(l*((a - xi*xip)^2 - 4*b*(xi + xip)), n);
d:= igcdex(4*xip^3 + 4*a*xip + 4*b, n, 'l');
if 1 < d then
    return d
end if;
xiip:= mod(l*((xip^2 - a)^2 - 8*b*xip), n)
end if;
[x, xii, xiip]
end proc

```

▼ 10.3. Projektív reprezentáció.

```

> #
# Doubling on an elliptic "curve" modulo n, where gcd(n,6)=1.
# x, z are the coordinates of the point to double and a, b
# are the parameters. The return value is a list, the
# coordinates of the double of the point.
#
ellxdou:=proc(x,z,a,b,n)
[(x^2-a*z^2)^2-8*b*x*z^3 mod n,4*z*(x^3+a*x*z^2+b*z^3) mod n]
end;
ellxdou:= proc(x, z, a, b, n)
[mod((x^2 - a*z^2)^2 - 8*b*x*z^3, n), mod(4*z*(x^3
+ a*x*z^2 + b*z^3), n)]
end proc
(10.3.1)
> #
# Calculation of the triple of coordinates
# [[x,z],[x_2i,z_2i],[x_2i+1,z_2i+1]] if t=0
# or [[x,z],[x_2i+1,z_2i+1],[x_2i+2,z_2i+2]] if t=1
# from the triple of coordinates [[x,z],[x_i,z_i],[x_i+1,
z_i+1]]
# on an elliptic "curve" modulo n, where igcd(n,6)=1, and a,
b
# are the parameters. The return value is this list.
#
ellnextxz:=proc(L,t,a,b,n) local l,d,x,z,xi,zi,xip,zip,xii,
zii,xiip,ziip;
x:=L[1][1]; z:=L[1][2];
xi:=L[2][1]; zi:=L[2][2]; xip:=L[3][1]; zip:=L[3][2];
if t=0 then
    xii:=(xi^2-a*zi^2)^2-8*b*xi*zi^3 mod n;

```

```

    zii:=4*zi*(xi^3+a*xi*zi^2+b*zi^3) mod n;
    xiip:=z*((xi*xip-a*zi*zip)^2-4*b*zi*zip*(xi*zip+xip*zi))
mod n;
    ziip:=x*(xip*zi-xi*zip)^2 mod n;
else
    xii:=z*((xi*xip-a*zi*zip)^2-4*b*zi*zip*(xi*zip+xip*zi)) mod
n;
    zii:=x*(xip*zi-xi*zip)^2 mod n;
    xiip:=(xip^2-a*zip^2)^2-8*b*xip*zip^3 mod n;
    ziip:=4*zip*(xip^3+a*xip*zip^2+b*zip^3) mod n;
fi; [[x,z],[xii,zii],[xiip,ziip]] end;
ellnextz:=proc(L,t,a,b,n)
local l,d,x,z,xi,zi,xip,zip,xii,zii,xiip,ziip;
x:=L[1][1];
z:=L[1][2];
xi:=L[2][1];
zi:=L[2][2];
xip:=L[3][1];
zip:=L[3][2];
if t = 0 then
    xii:=mod((xi^2-a*zi^2)^2-8*b*xi*zi^3,n);
    zii:=mod(4*zi*(xi^3+a*xi*zi^2+b*zi^3),n);
    xiip:=mod(z*((xi*xip-a*zi*zip)^2-4*b*zi*zip*(xi*zip
+xip*zi)),n);
    ziip:=mod(x*(xip*zi-xi*zip)^2,n)
else
    xii:=mod(z*((xi*xip-a*zi*zip)^2-4*b*zi*zip*(xi*zip
+xip*zi)),n);
    zii:=mod(x*(xip*zi-xi*zip)^2,n);
    xiip:=mod((xip^2-a*zip^2)^2-8*b*xip*zip^3,n);
    ziip:=mod(4*zip*(xip^3+a*xip*zip^2+b*zip^3),n)
end if;
[[x,z],[xii,zii],[xiip,ziip]]
end proc
> #
# This program compute k*P, k>=0 or a divisor of n
# for a list of elliptic "curve" modulo n, where gcd(n,6)=1.
# It use the left-to-right binary method,
# projective representation, and backtracking.
#
batchellmults:=proc(L,k,n) local LL,LLL,LLLL,a,b,i,j,P,B,d,l;

```

(10.3.2)

```

if k=0 then return L fi; LL:=[];
for i to nops(L) do
  a:=L[i][2]; b:=L[i][3];
  P:=ellxzdou(L[i][1],1,a,b,n);
  P:=[[L[i][1],1],[L[i][1],1],P,L[i][2],L[i][3]];
  LL:=[op(LL),P];
od;
B:=convert(k,base,2);
for j from nops(B)-1 to 1 by -1 do
  LLL:=LL; LL:=[];
  for i to nops(LLL) do
    P:=LLL[i][1..3]; a:=LLL[i][4]; b:=LLL[i][5];
    P:=ellnextxz(P,B[j],a,b,n);
    LL:=[op(LL),[op(P),a,b]];
  od;
od; LLL:=[];
for i to nops(LL) do LLL:=[op(LLL),LL[i][2][2]] od;
LLL:=batchmodinv(LLL,n);
if type(LLL,integer) then
  if LLL<n then return LLL fi; LL:=[];
  for i to nops(L) do
    a:=L[i][2]; b:=L[i][3];
    P:=ellxzdou(L[i][1],1,a,b,n);
    d:=igcdex(P[2],n,'1');
    if 1<d then
      if d<n then return d else next fi;
    else
      P:=[[L[i][1],1],[L[i][1],1],[1*P[1] mod n,1]];
      fi;
      for j from nops(B)-1 to 1 by -1 do
        P:=ellnextxz(P,B[j],a,b,n);
        d:=igcdex(P[2][2],n,'1');
        if 1<d then
          if d<n then return d else break fi;
        else
          P[2][1]:=1*P[2][1]*1 mod n; P[2][2]:=1;
          fi;
          d:=igcdex(P[2][2],n,'1');
          if 1<d then
            if d<n then return d else break fi;
          else
            P[3][1]:=1*P[3][1] mod n; P[3][2]:=1;
            fi;
          od;
          if d=n then next fi;
          LL:=[op(LL),[P[2][1],a,b]];
        od;
      LLLL:=LL;

```

```

else
  LLLL:=[];
  for i to nops(LLL) do
    P:=LL[i][2][1]*LLL[i] mod n;
    a:=LL[i][4]; b:=LL[i][5];
    LLLL:=[op(LLLL), [P,a,b]];
  od;
fi; LLLL; end;
batchllmults:= proc(L, k, n)
local LL, LLL, LLLL, a, b, i, j, P, B, d, t;
if k = 0 then
  return L
end if;
LL:= [];
for ito nops(L) do
  a:= L[i][2];
  b:= L[i][3];
  P:= ellxzdou(L[i][1], 1, a, b, n);
  P:= [[L[i][1], 1], [L[i][1], 1], P, L[i][2], L[i][3]];
  LL:= [op(LL), P]
end do;
B:= convert(k, base, 2);
for j from nops(B) - 1 by -1 to 1 do
  LLL:= LL;
  LL:= [];
  for ito nops(LLL) do
    P:= LLL[i][1..3];
    a:= LLL[i][4];
    b:= LLL[i][5];
    P:= ellnextxz(P, B[j], a, b, n);
    LL:= [op(LL), [op(P), a, b]]
  end do
end do;
LLL:= [];
for ito nops(LL) do
  LLL:= [op(LLL), LL[i][2][2]]
end do;
LLL:= batchmodinv(LLL, n);
if type(LLL, integer) then

```

(10.3.3)

```

if  $LLL < n$  then
    return  $LLL$ 
end if;
 $LL := []$ ;
for  $i$  to  $nops(L)$  do
     $a := L[i][2]$ ;
     $b := L[i][3]$ ;
     $P := ellxzdou(L[i][1], 1, a, b, n)$ ;
     $d := igcdex(P[2], n, 'T')$ ;
    if  $1 < d$  then
        if  $d < n$  then
            return  $d$ 
        else
            next
        end if
    else
         $P := [[L[i][1], 1], [L[i][1], 1], [mod(l * P[1], n), 1]]$ 
    end if;
    for  $j$  from  $nops(B) - 1$  by  $-1$  to  $1$  do
         $P := ellnextxz(P, B[j], a, b, n)$ ;
         $d := igcdex(P[2][2], n, 'T')$ ;
        if  $1 < d$  then
            if  $d < n$  then
                return  $d$ 
            else
                break
            end if
        else
             $P[2][1] := mod(l * P[2][1] * l, n)$ ;
             $P[2][2] := 1$ 
        end if;
         $d := igcdex(P[2][2], n, 'T')$ ;
        if  $1 < d$  then
            if  $d < n$  then
                return  $d$ 
            else
                break
            end if

```

```

        else
            P[3][1] := mod(l*P[3][1], n);
            P[3][2] := 1
        end if
    end do;
    if d = n then
        next
    end if;
    LL := [op(LL), [P[2][1], a, b]]
end do;
LLLL := LL
else
    LLLL := [];
    for i to nops(LLL) do
        P := mod(LL[i][2][1]*LLL[i], n);
        a := LL[i][4];
        b := LL[i][5];
        LLLL := [op(LLLL), [P, a, b]]
    end do
end if;
LLLL
end proc
> #
# This procedure is elliptic curve method with projective
# representation for factorization of n. It use K "curves"
# parallel and only after all step of a multiplication is
# done an igcd operation. If no success,
# backtracking is used for each curves separately.
Unsuccessful
# curves are "killed" from the list of curves.
# The prime powers up to P are considered so that they
# are not less then the bound B. The result is the list
# of triples [x,a,b], where x is the first coordinate of
# the multiple and a,b are the parameters of the curve.
#
batchllsplit:=proc(n,K,P,B) local e,d,p,L,i;
L:=[];
for i to K do
    p:=ellrand(n); if type(p,integer) then return p fi;
    L:=[op(L), [p[1],p[3],p[4]]];
od;

```



```

p:=2; e:=1; while 2^e<B do e:=e+1 od;
while p<=P do
  while p^e>p*B and e>1 do e:=e-1 od;
  L:=batchellmuls(L,p^e,n);
  if type(L,integer) then return L fi;
  p:=nextprime(p);
od; L; end;
batchellsplit := proc(n, K, P, B)
  local e, d, p, L, i;
  L := [ ];
  for i to K do
    p := ellrand(n);
    if type(p, integer) then
      return p
    end if;
    L := [op(L), [p[1], p[3], p[4]]]
  end do;
  p := 2;
  e := 1;
  while 2^e < B do
    e := e + 1
  end do;
  while p <= P do
    while p*B < p^e and 1 < e do
      e := e - 1
    end do;
    L := batchellmuls(L, p^e, n);
    if type(L, integer) then
      return L
    end if;
    p := nextprime(p)
  end do;
  L
end proc

```

(10.3.4)

>

▼ 10.4. Második lépcső.

This is a calculation of the distribution function F of the size of the largest

prime factor of a random number.

> **F2:=x->1+ln(x);**

$$F2 := x \rightarrow 1 + \ln(x) \quad (10.4.1)$$

> **F3:=x->F2(1/2)-int(F2(t/(1-t))/t,t=x..1/2);**

$$F3 := x \rightarrow F2\left(\frac{1}{2}\right) - \int_x^{\frac{1}{2}} \frac{F2\left(\frac{t}{1-t}\right)}{t} dt \quad (10.4.2)$$

> **F4:=x->F3(1/3)-int(F3(t/(1-t))/t,t=x..1/3);**

$$F4 := x \rightarrow F3\left(\frac{1}{3}\right) - \int_x^{\frac{1}{3}} \frac{F3\left(\frac{t}{1-t}\right)}{t} dt \quad (10.4.3)$$

The next line would be the definition between 1/5 and 1/4, but too slow; instead we will use approximation.

> **F5:=x->F4(1/4)-int(F4(t/(1-t))/t,t=x..1/4);**

$$F5 := x \rightarrow F4\left(\frac{1}{4}\right) - \int_x^{\frac{1}{4}} \frac{F4\left(\frac{t}{1-t}\right)}{t} dt \quad (10.4.4)$$

> **evalf(F4(0.25));**

$$0.0049109254 \quad (10.4.5)$$

> **Order:=10; series(F4(x),x=1/4);**

Order:= 10

$$1 - \ln(3) - \frac{1}{2} \ln(2)^2 - \operatorname{dilog}\left(\frac{3}{2}\right) \quad (10.4.6)$$

$$\begin{aligned}
& + \ln(2) \ln(3) - \frac{1}{12} \pi^2 - \int_{\frac{1}{4}}^{\frac{1}{3}} \frac{1 - \ln(2) - \int_{\frac{t}{1-t}}^{\frac{1}{2}} \frac{1 + \ln\left(\frac{t}{1-t}\right)}{t} dt}{t} dt \\
& + \left(4 - 4 \ln(3) - 2 \ln(2)^2 - 4 \operatorname{dilog}\left(\frac{3}{2}\right) \right) \\
& + 4 \ln(2) \ln(3) - \frac{1}{3} \pi^2 \left(x - \frac{1}{4} \right) + \left(\frac{8}{3} - \frac{32}{3} \ln(2) + 8 \ln(3) \right) \\
& + 4 \ln(2)^2 + 8 \operatorname{dilog}\left(\frac{3}{2}\right) - 8 \ln(2) \ln(3) + \frac{2}{3} \pi^2 \left(x - \frac{1}{4} \right)^2 + \left(\frac{320}{27} \right. \\
& + \frac{1024}{27} \ln(2) - \frac{64}{3} \ln(3) - \frac{32}{3} \ln(2)^2 - \frac{64}{3} \operatorname{dilog}\left(\frac{3}{2}\right) \\
& + \left. \frac{64}{3} \ln(2) \ln(3) - \frac{16}{9} \pi^2 \right) \left(x - \frac{1}{4} \right)^3 + \left(-\frac{4160}{81} - \frac{11008}{81} \ln(2) \right) \\
& + 64 \ln(3) + 32 \ln(2)^2 + 64 \operatorname{dilog}\left(\frac{3}{2}\right) - 64 \ln(2) \ln(3) \\
& + \frac{16}{3} \pi^2 \left(x - \frac{1}{4} \right)^4 + \left(\frac{65536}{135} \ln(2) \right) \\
& + \frac{7168}{27} - \frac{1024}{5} \ln(3) - \frac{512}{5} \ln(2)^2 - \frac{1024}{5} \operatorname{dilog}\left(\frac{3}{2}\right) \\
& + \frac{1024}{5} \ln(2) \ln(3) - \frac{256}{15} \pi^2 \left(x - \frac{1}{4} \right)^5 \\
& + \left(-\frac{800768}{729} - \frac{6397952}{3645} \ln(2) + \frac{2048}{3} \ln(3) + \frac{1024}{3} \ln(2)^2 \right. \\
& + \left. \frac{2048}{3} \operatorname{dilog}\left(\frac{3}{2}\right) - \frac{2048}{3} \ln(2) \ln(3) + \frac{512}{9} \pi^2 \right) \left(x - \frac{1}{4} \right)^6 \\
& + \left(\frac{354992128}{76545} + \frac{490471424}{76545} \ln(2) - \frac{16384}{7} \ln(3) - \frac{8192}{7} \ln(2)^2 \right. \\
& - \left. \frac{16384}{7} \operatorname{dilog}\left(\frac{3}{2}\right) + \frac{16384}{7} \ln(2) \ln(3) - \frac{4096}{21} \pi^2 \right) \left(x - \frac{1}{4} \right)^7
\end{aligned}$$

$$\begin{aligned}
& + \left(-\frac{1422598144}{76545} - \frac{1806270464}{76545} \ln(2) + 8192 \ln(3) + 4096 \ln(2)^2 \right. \\
& + 8192 \operatorname{dilog}\left(\frac{3}{2}\right) - 8192 \ln(2) \ln(3) + \frac{2048}{3} \pi^2 \left. \right) \left(x - \frac{1}{4}\right)^8 \\
& + \left(\frac{155127119872}{2066715} \right. \\
& + \frac{180925497344}{2066715} \ln(2) - \frac{262144}{9} \ln(3) - \frac{131072}{9} \ln(2)^2 \\
& \left. - \frac{262144}{9} \operatorname{dilog}\left(\frac{3}{2}\right) + \frac{262144}{9} \ln(2) \ln(3) - \frac{65536}{27} \pi^2 \right) \left(x - \frac{1}{4}\right)^9 \\
& + O\left(\left(x - \frac{1}{4}\right)^{10}\right)
\end{aligned}$$

> F4s:=evalf(%);

$$\begin{aligned}
F4s := & 0.00491092536 + 0.194433553 (x - 0.2500000000) & (10.4.7) \\
& + 2.884229635 (x - 0.2500000000)^2 \\
& + 17.84374609 (x - 0.2500000000)^3 \\
& - 84.66851915 (x - 0.2500000000)^4 \\
& + 407.1260619 (x - 0.2500000000)^5 \\
& - 1665.622940 (x - 0.2500000000)^6 \\
& + 6852.317002 (x - 0.2500000000)^7 \\
& - 27147.86069 (x - 0.2500000000)^8 \\
& + 1.080283354 \cdot 10^5 (x - 0.2500000000)^9 \\
& + O((x - 0.2500000000)^{10})
\end{aligned}$$

> F4p:=convert(F4s,polynomial); F4as:=subs(x=t/(1-t),F4p)/t;

$$\begin{aligned}
F4p := & -0.04369746289 + 0.194433553 x \\
& + 2.884229635 (x - 0.2500000000)^2 \\
& + 17.84374609 (x - 0.2500000000)^3 \\
& - 84.66851915 (x - 0.2500000000)^4 \\
& + 407.1260619 (x - 0.2500000000)^5 \\
& - 1665.622940 (x - 0.2500000000)^6 \\
& + 6852.317002 (x - 0.2500000000)^7 \\
& - 27147.86069 (x - 0.2500000000)^8
\end{aligned}$$

$$\begin{aligned}
& + 1.080283354 \cdot 10^5 (x - 0.2500000000)^9 \\
F4as := & \frac{1}{t} \left(-0.04369746289 + \frac{0.194433553 t}{1-t} \right. \\
& + 2.884229635 \left(\frac{t}{1-t} - 0.2500000000 \right)^2 \\
& + 17.84374609 \left(\frac{t}{1-t} - 0.2500000000 \right)^3 \\
& - 84.66851915 \left(\frac{t}{1-t} - 0.2500000000 \right)^4 \\
& + 407.1260619 \left(\frac{t}{1-t} - 0.2500000000 \right)^5 \\
& - 1665.622940 \left(\frac{t}{1-t} - 0.2500000000 \right)^6 \\
& + 6852.317002 \left(\frac{t}{1-t} - 0.2500000000 \right)^7 \\
& - 27147.86069 \left(\frac{t}{1-t} - 0.2500000000 \right)^8 \\
& \left. + 1.080283354 \cdot 10^5 \left(\frac{t}{1-t} - 0.2500000000 \right)^9 \right)
\end{aligned} \tag{10.4.8}$$

> **F5a:=y->F4(1/4)-int(F4as,t=y..1/4);**

$$F5a := y \rightarrow F4\left(\frac{1}{4}\right) - \int_y^{\frac{1}{4}} F4as \, dt \tag{10.4.9}$$

> **evalf(F5a(0.20));**

$$0.00035463226 \tag{10.4.10}$$

> **F:=proc(x) if x>1. then 1. elif x>=1/2. then F2(x) elif x>=1/3. then F3(x) elif x>=1/4. then F4(x) elif x>=1/5. then F5a(x) else 0. fi; end;**

F:=proc(x) (10.4.11)

if 1. < x then

1.

elif 1 / 2. <= x then

$F2(x)$

elif 1 / 3. <= x then

$F3(x)$

elif 1 / 4. <= x **then**

$F_4(x)$

elif 1 / 5. <= x **then**

$F_{5a}(x)$

else

0.

end if

end proc

> **F(0.22);**

$$0.9963042122 - \ln(3) - \frac{1}{2} \ln(2)^2 - \operatorname{dilog}\left(\frac{3}{2}\right) \quad (10.4.12)$$

$$+ \ln(2) \ln(3) - \frac{1}{12} \pi^2 - \int_{\frac{1}{4}}^{\frac{1}{3}} \frac{1 - \ln(2) - \int_{\frac{t}{1-t}}^{\frac{1}{2}} \frac{1 + \ln\left(\frac{t}{1-t}\right)}{t} dt}{t} dt$$

> **evalf(%);**

0.00121513756

(10.4.13)

Now the definition of F is finished. We start to count probability ENI and ENII to does not succed with working N elliptic curves, and using step I or step II also. We have to give the following bounds:

B0 is the bound for prime powers; B1 is the bound for primes in step I; B2 is the bound for primes in step II; finally, B is the bound for factors; the conditions $B_0 \geq B_1$ and $B_2 > B_1$ are supposed.

The probabilities to does not succeed with one elliptic curve is also given, these are EI and EII, respectively. Two additional probabilities also given:

E0 is an upper bound for probability that error is caused by too small B0, and E1 is an upper bound the probability that error is cause by too small B0 with respect to B1. These shoud have to keep below 1-EI or 1-EII, depending on whether we plane to use step I only or step II, too.

Upper bounds for work also given back: these are calculated for one curve as the number of additions on the curve, and are the following:

W0 for part of step I depending only on B0;

W1 for part of step I depending mainly on B1;

W2 for step II and depending mainly on B2.

```

> with(numtheory);
[GIgcd, bigomega, cfrac, cfracpol, cyclotomic, divisors, factorEQ,
factorset, fermat, imagunit, index, integral_basis, invcfrac, invphi,
issqrfree, jacobi, kronecker,  $\lambda$ , legendre, mcombine, mersenne,
migcdex, minkowski, mipolys, mlog, mobius, mroot, msqrt, nearestp,
nthconver, nthdenom, nthnumer, nthpow, order, pdexpand,  $\phi$ ,  $\pi$ ,
pprimroot, primroot, quadres, rootsunity, safeprime,  $\sigma$ , sq2factor,
sum2sqr,  $\tau$ , thue]

```

(10.4.14)

```

> B0:=10.^5; B1:=10.^4; B2:=10.^6; B:=10.^20; N:=32;
      B0:= 1.00000 105
      B1:= 10000.
      B2:= 1.000000 106
      B:= 1.000000000 1020
      N:= 32

```

(10.4.15)

```

> E0:=proc() evalf(pi(floor(sqrt(B0)))/B0); end; E0();
E0:= proc() evalf(numtheory:- $\pi$ (floor(sqrt(B0)))/B0) end proc
      0.0006500000000

```

(10.4.16)

```

> W0:=proc() 2*floor(log[2.](B0))*pi(floor(sqrt(B0)))*N; end;
W0();
W0:= proc()
      2*floor(log[2.](B0))*numtheory:- $\pi$ (floor(sqrt(B0)))*
      N
end proc
      66560

```

(10.4.17)

```

> E1:=proc() local s,p; p:=nextprime(floor(sqrt(B0))); s:=0.;
while p<B1 do s:=s+1/p^2; p:=nextprime(p); od; s; end;
E1:= proc()

```

(10.4.18)

```

      local s, p;
      p:= nextprime(floor(sqrt(B0)));
      s:= 0.;
      while p < B1 do
          s:= s + 1 / p^2;
          p:= nextprime(p)
      end do;
      s
end proc

```

```

> E1();
      0.0004499193528

```

(10.4.19)

```

> W1:=proc() local s,p; p:=nextprime(floor(sqrt(B0))); s:=0;
  while p<B1 do s:=s+floor(log[2.](p)); p:=nextprime(p); od; 2*
  N*s; end;
W1:=proc()

```

(10.4.20)

```

  local s, p;
  p:= nextprime(floor(sqrt(B0)));
  s:= 0;
  while p < B1 do
    s:= s + floor(log[2.](p));
    p:= nextprime(p)
  end do;
  2 * N * s

```

```
end proc
```

```

> W1();

```

850368 (10.4.21)

```

> EI:=proc() evalf(1.-F(log[B](B1))); end;
  EI(); 1-EI(); ENI:=proc() EI()^N; end; ENI();
  EI:=proc() evalf(1. - F(log[B](B1))) end proc

```

0.9996453676

0.0003546324

ENI:=proc() EI()^N end proc

0.9887139216

(10.4.22)

```

> EII:=proc() evalf(1.-exp(gamma)*log[B](B1)*F(log[B](B2)));
  end;

```

```

  EII(); 1-EII(); ENII:=proc() EII()^N; end; ENII();

```

```

EII:=proc() evalf(1. - exp( $\gamma$ ) * log[B](B1) * F(log[B](B2))) end proc

```

0.9915754783

0.0084245217

ENII:=proc() EII()^N end proc

0.7628249680

(10.4.23)

```

> W2:=proc() ceil(pi(floor(B2))-pi(floor(B1)))*N; end; W2();

```

```
W2:=proc()
```

```
  ceil(numtheory:- $\pi$ (floor(B2))
```

```
  - numtheory:- $\pi$ (floor(B1)))*N
```

```
end proc
```

2472608

(10.4.24)

```

> N:=64; B0:=10.^6; B1:=10.^5; B2:=10.^7; B:=10.^25;
  E0(); E1(); EI(); 1-EI(); ENI(); EII(); 1-EII(); ENII();
  W0(); W1(); W2();

```

N:= 64


```

BO:= 1.000000 106
BI:= 1.000000 105
B2:= 1.0000000 107
B:= 1.000000000 1025
0.0001680000000
0.0001261875417
0.9996453676
0.0003546324
0.9775552187
0.9950978649
0.0049021351
0.7301479030
408576
17655424
41919168

```

(10.4.25)

```

> #
# This procedure is the second step of elliptic curve method
# for
# factorization. The list of "curves" is L, the table of
# doubles
# has size N and primes greater than m but not greater than M
# are considered. The result is the last list or a factor of
# n.
#

batchellsplit2:=proc(L,n::posint,N::posint,m::posint,
M::posint)
local y,i,j,E,P,PP,p,pp,d,LL,LLL,a,b;
LL:=[];
for i to nops(L) do
p:=L[i];
y:=msqrt(L[i][1],n);
if y=FAIL then next fi;
LL:=[op(LL),[L[i][1],y,1,L[i][2],L[i][3]]];
od;
E:=Array(1..nops(LL),1..N);
for i to nops(LL) do
P:=LL[i][1..3];
a:=LL[i][4]; b:=LL[i][5];
PP:=elldou(P,a,b,n);
if type(PP,integer) then return PP fi;
E[i,1]:=PP;
for j from 2 to N do

```

```

        PP:=elladd(E[i,j-1],PP,a,b,n);
        if type(PP,integer) then return PP fi;
        E[i,1]:=PP;
    od;
od;
p:=nextprime(m);
for i to nops(LL) do
    P:=LL[i][1..3];
    a:=LL[i][4]; b:=LL[i][5];
    PP:=ellmul(P,p,a,b,n);
    if type(PP,integer) then return PP fi;
    LLL[i]:=[op(PP),a,b];
od;
pp:=nextprime(p);
while pp<M do
d:=pp-p; p:=pp;
    if d<=2*N then
        for i to nops(LL) do
            PP:=LLL[i][1..3];
            a:=LL[i][4]; b:=LL[i][5];
            PP:=elladd(PP,E[i,d/2],a,b,n);
            if type(PP,integer) then return PP fi;
            LLL[i]:=[op(PP),a,b];
        od;
    else
        for i to nops(LL) do
            P:=LL[i][1..3];
            a:=LL[i][4]; b:=LL[i][5];
            PP:=ellmul(P,d,a,b,n);
            if type(PP,integer) then return PP fi;
            PP:=elladd(PP,LLL[i][1..3],a,b,n);
            if type(PP,integer) then return PP fi;
            LLL[i]:=[op(PP),a,b];
        od;
    fi;
pp:=nextprime(p);
od; LLL; end;

```

batchsplit2 := **proc**(*L*, *n*::*posint*, *N*::*posint*, *m*::*posint*, *M*::*posint*)

(10.4.26)

local *y*, *i*, *j*, *E*, *P*, *PP*, *p*, *pp*, *d*, *LL*, *LLL*, *a*, *b*;

LL := [];

for *i* to *nops*(*L*) **do**

p := *L*[*i*];

y := *msqrt*(*L*[*i*][1], *n*);

if *y* = *FAIL* **then**

next

end if;

```

    LL := [op(LL), [L[i][1], y, 1, L[i][2], L[i][3]]]
end do;
E := Array(1..nops(LL), 1..N);
for ito nops(LL) do
    P := LL[i][1..3];
    a := LL[i][4];
    b := LL[i][5];
    PP := elldou(P, a, b, n);
    if type(PP, integer) then
        return PP
    end if;
    E[i, 1] := PP;
    for j from 2 to N do
        PP := elladd(E[i, j - 1], PP, a, b, n);
        if type(PP, integer) then
            return PP
        end if;
        E[i, 1] := PP
    end do
end do;
p := nextprime(m);
for ito nops(LL) do
    P := LL[i][1..3];
    a := LL[i][4];
    b := LL[i][5];
    PP := ellmul(P, p, a, b, n);
    if type(PP, integer) then
        return PP
    end if;
    LLL[i] := [op(PP), a, b]
end do;
pp := nextprime(p);
while pp < M do
    d := pp - p;
    p := pp;
    if d <= 2 * N then
        for ito nops(LL) do
            PP := LLL[i][1..3];

```

```

    a := LL[i][4];
    b := LL[i][5];
    PP := elladd(PP, E[i, 1 / 2 * d], a, b, n);
    if type(PP, integer) then
        return PP
    end if;
    LLL[i] := [op(PP), a, b]
end do
else
    for ito nops(LL) do
        P := LL[i][1..3];
        a := LL[i][4];
        b := LL[i][5];
        PP := ellmul(P, d, a, b, n);
        if type(PP, integer) then
            return PP
        end if;
        PP := elladd(PP, LLL[i][1..3], a, b, n);
        if type(PP, integer) then
            return PP
        end if;
        LLL[i] := [op(PP), a, b]
    end do
end if;
pp := nextprime(p)
end do;
LLL
end proc

```

- ▶ 11. Prímteszt elliptikus görbékkel
- ▶ 12. Polinomfaktorizálás
- ▶ 13. Az AKS-teszt
- ▶ 14. A szita módszerek alapjai

- ▶ **15. Számtest szita**
- ▶ **16. Vegyes problémák**